

Dynamical Approach for Solving Complex Networks

Sidney Redner, Boston University, physics.bu.edu/~redner
collaborators: P. Krapivsky, F. Leyvraz

MECO34: 34th Conference of the Middle European Cooperation in Statistical Physics
30 March - 01 April 2009 Universität Leipzig, Germany

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Outline:

preliminaries: 1. Erdős-Rényi Random Graph
2. Random Recursive Tree (RRT)

linear preferential attachment by *redirection*

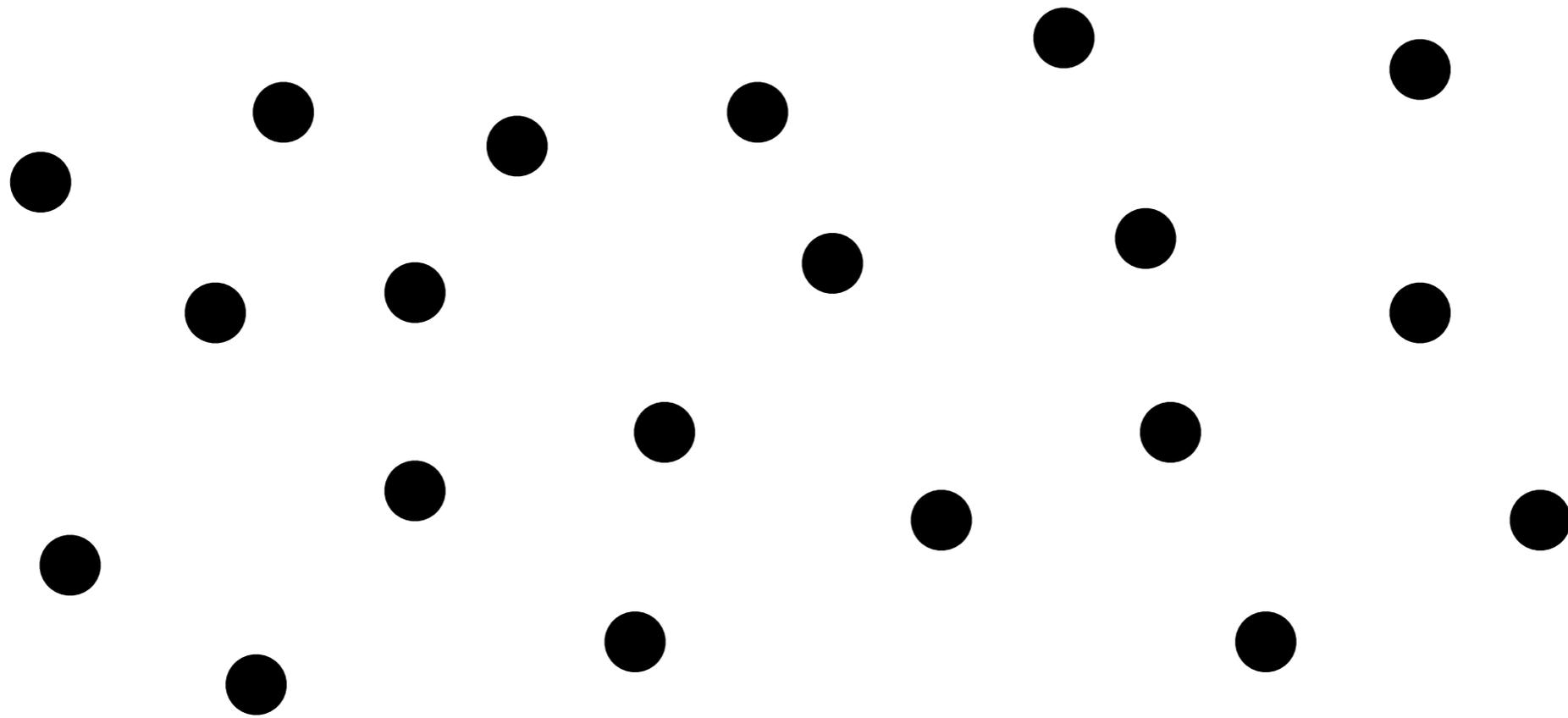
universality classes of preferential attachment

genealogy

outlook

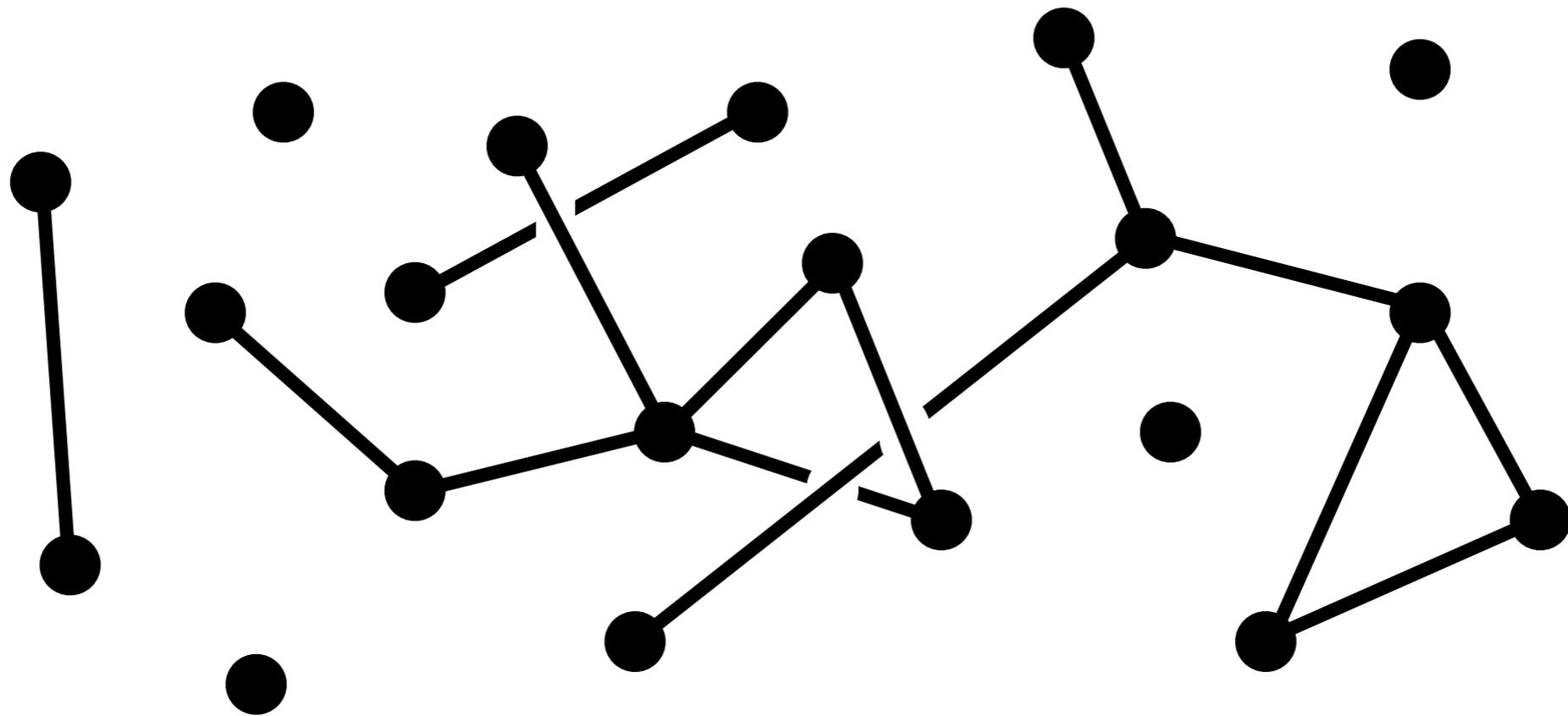
Erdős-Rényi Random Graph

Erdős-Rényi Random Graph



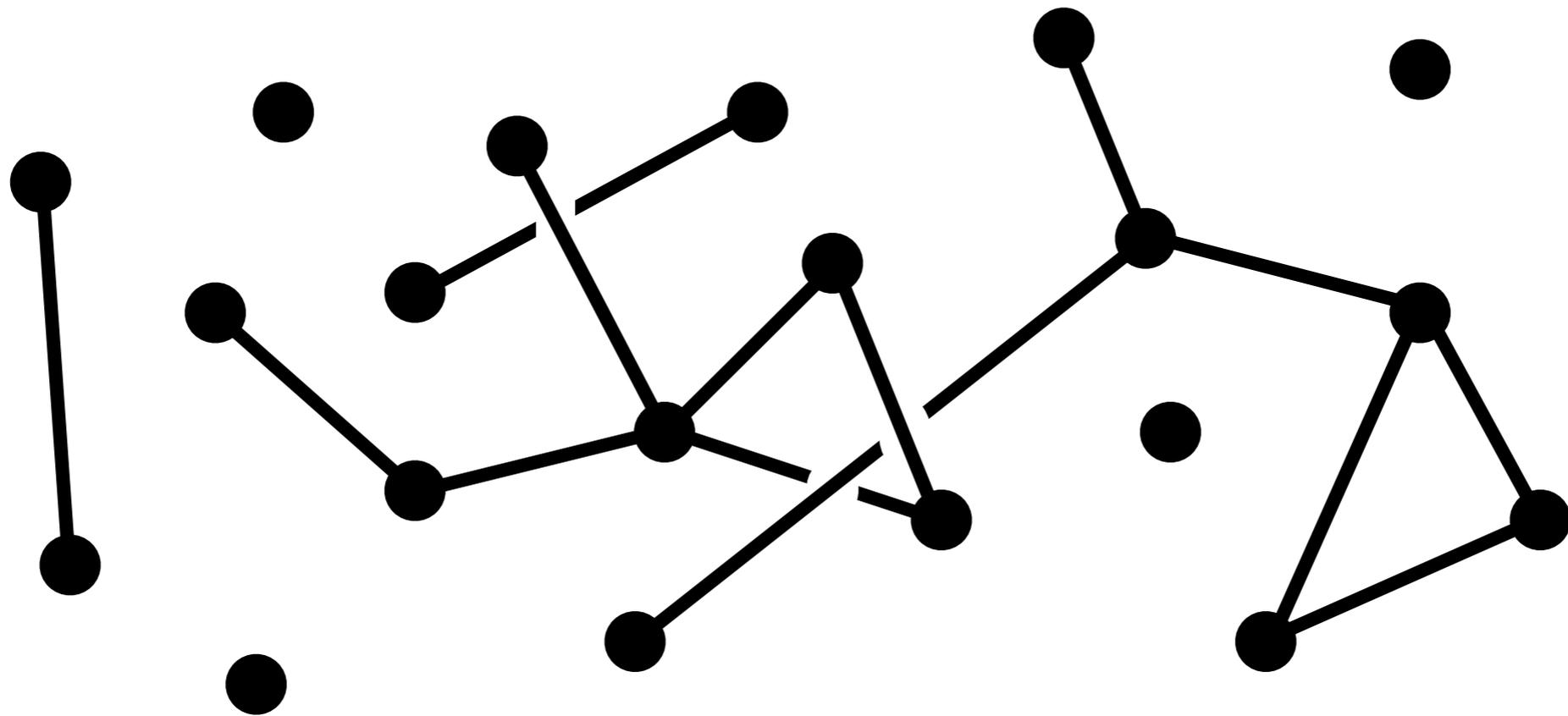
I. Start with N isolated nodes (and $N^2/2$ pairs)

Erdős-Rényi Random Graph



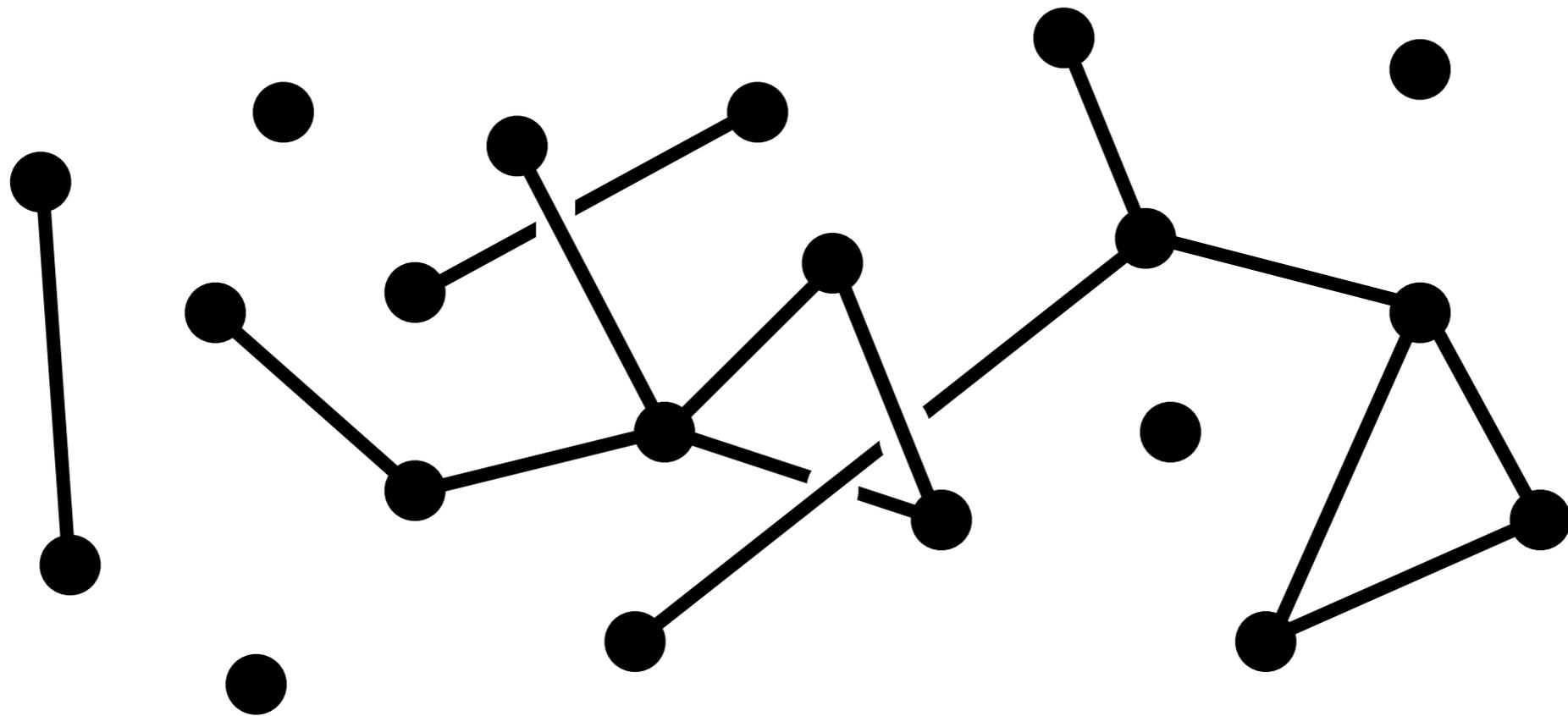
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2. Introduce links at constant rate

Erdős-Rényi Random Graph



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2. Introduce links at constant rate $N/2$

Erdős-Rényi Random Graph

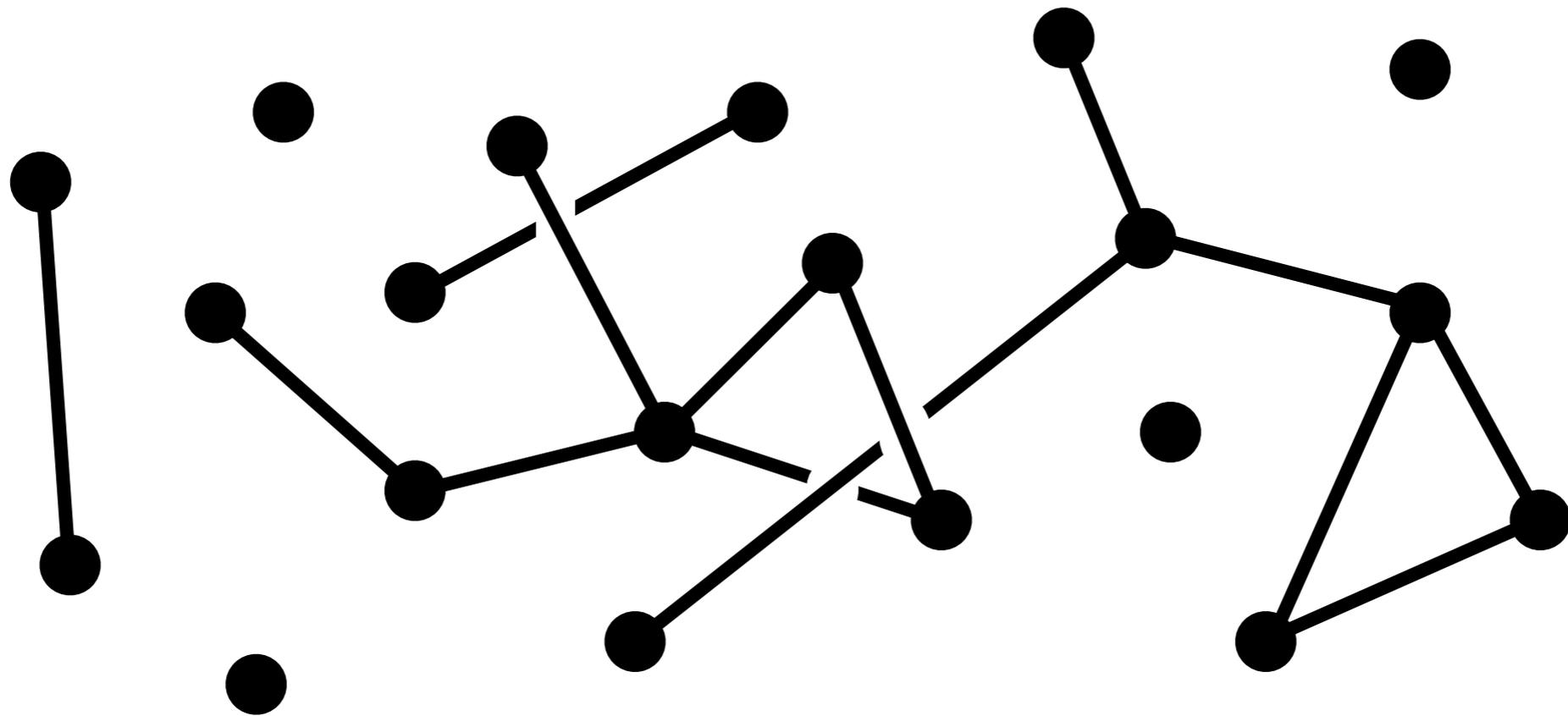


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Number of links L at time t : $L = N/2 \times t = Nt/2$

Average degree $D = 2L/N = t$

Erdős-Rényi Random Graph



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2. Introduce links at constant rate $N/2$

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Average degree $D = 2L/N = t$

Average degree $k \rightarrow k+1$ at rate 1

The Degree Distribution

N_k = number of nodes of degree k

n_k = fraction of nodes of degree k

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Master equation for n_k

$$\frac{dn_k}{dt} = n_{k-1} - n_k$$

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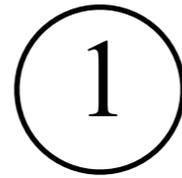
Solution

$$n_k = \frac{t^k}{k!} e^{-kt}$$

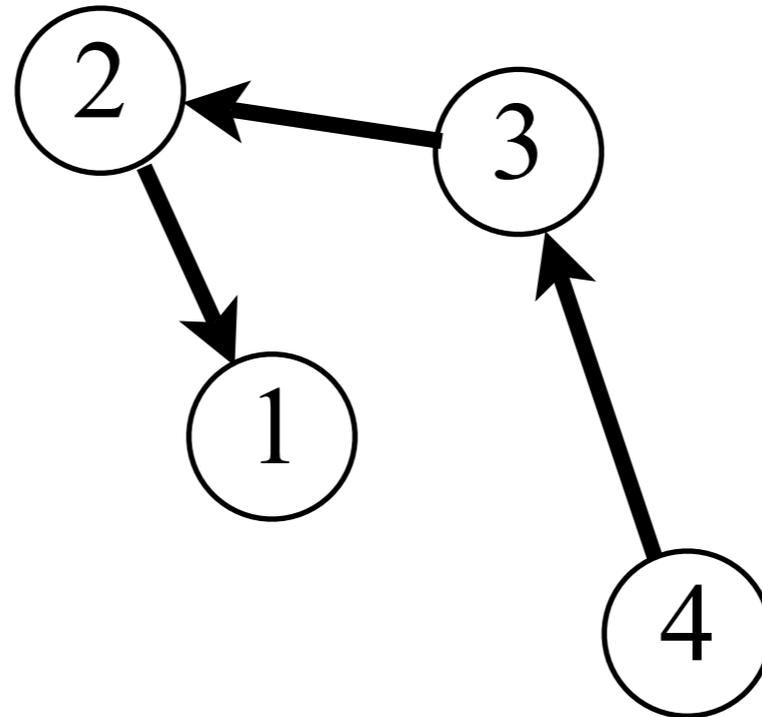
Poisson degree distribution

Random Recursive Tree

Random Recursive Tree

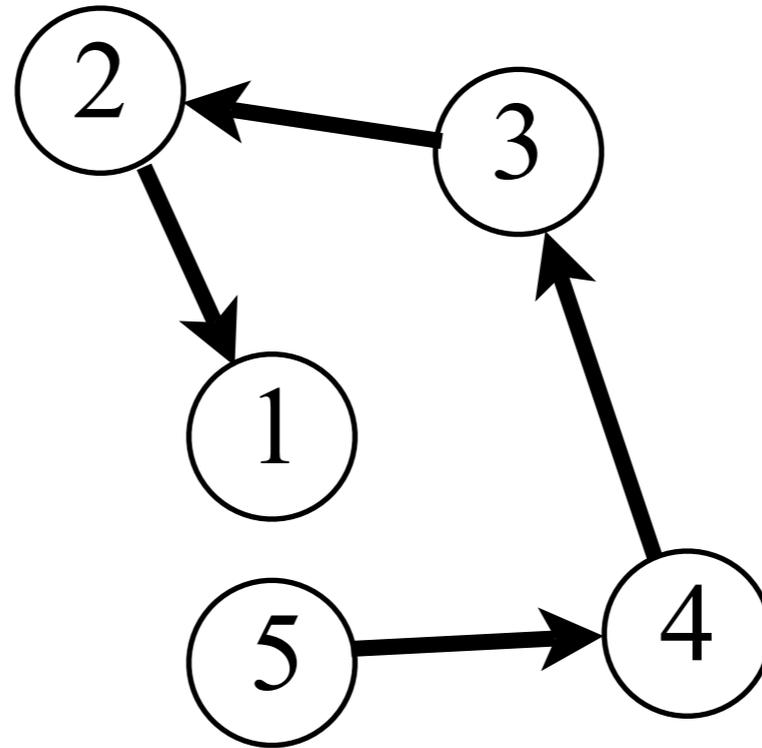


Random Recursive Tree



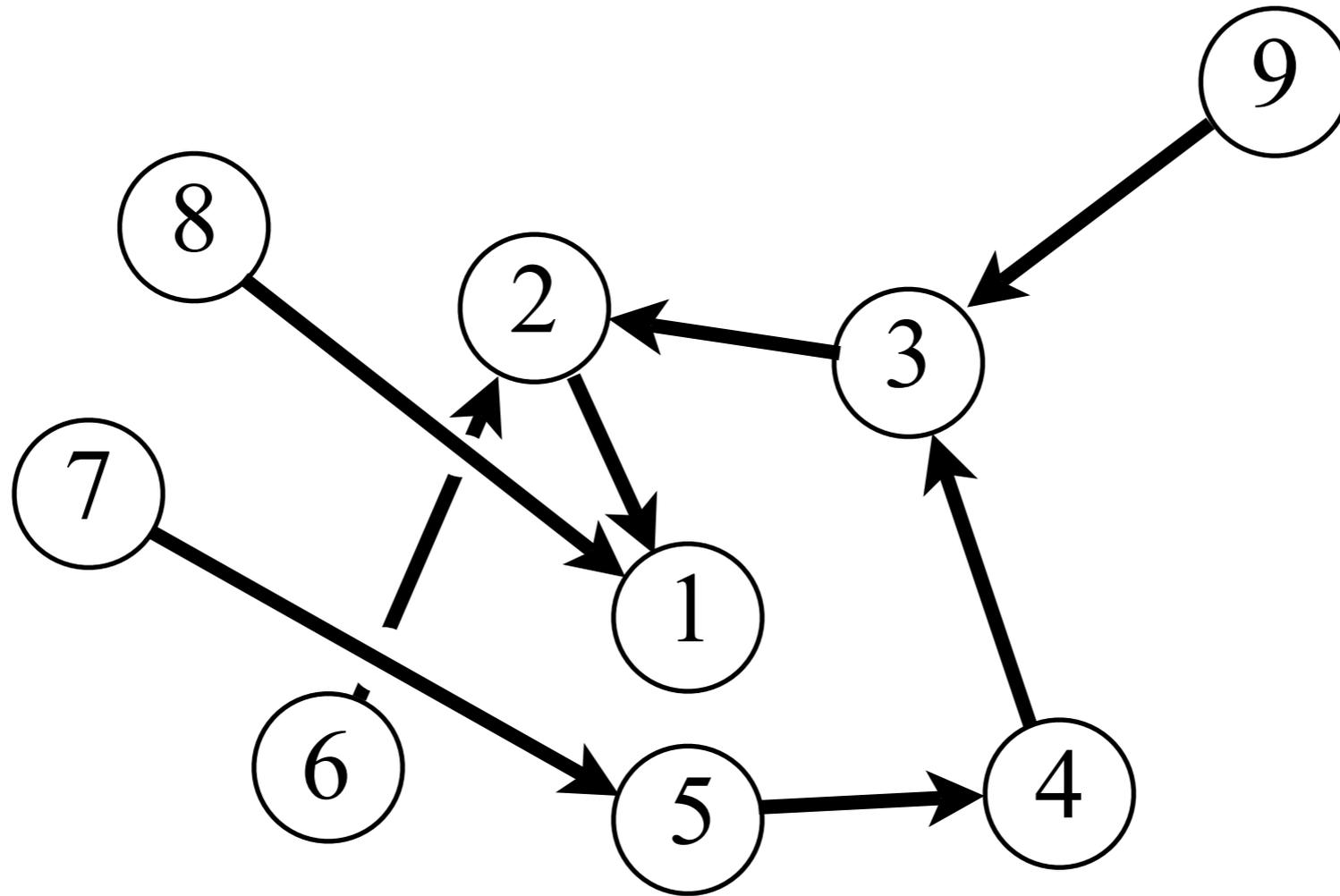
I. Introduce nodes one at a time

Random Recursive Tree



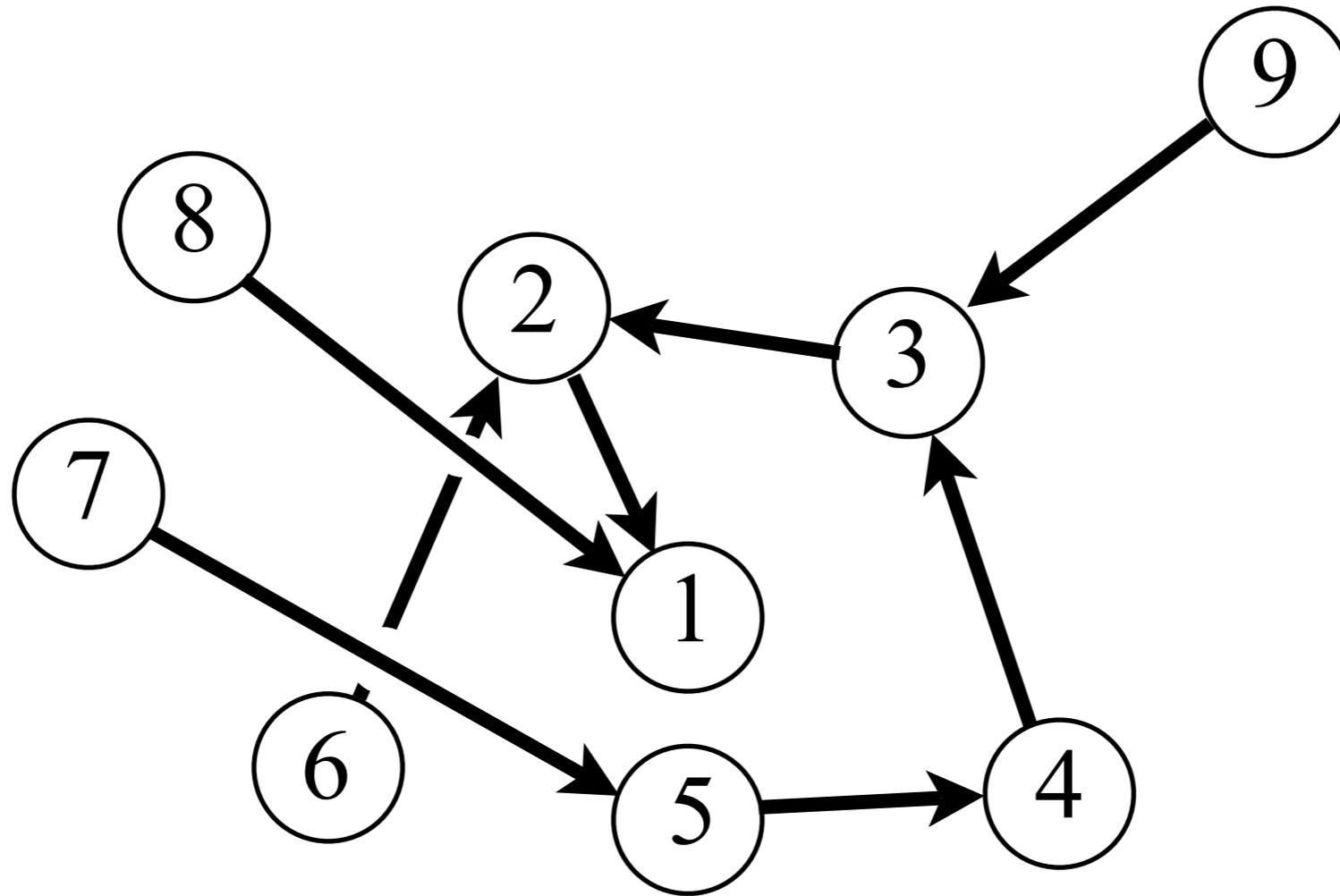
1. Introduce nodes one at a time
2. Attach to one earlier node *randomly and uniformly*

Random Recursive Tree



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Random Recursive Tree



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Basic observable: N_k , the degree distribution

The Degree Distribution

The Degree Distribution

Master
Equation

$$\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1}$$

The Degree Distribution

Master
Equation

$$\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1} \quad \frac{dn_k}{dt} = n_{k-1} - n_k$$

The Degree Distribution

Master
Equation

$$\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1}$$

“time”
variable N

$$\frac{dn_k}{dt} = n_{k-1} - n_k$$

The Degree Distribution

Master
Equation

$$\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1}$$

$\frac{dn_k}{dt} = n_{k-1} - n_k$

“time” variable N

gain of nodes of degree k

loss of nodes of degree k

The diagram shows the Master Equation for the degree distribution of nodes in a network. The equation is $\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1}$. The term $\frac{dN_k}{dN}$ is annotated with a blue arrow pointing to it from the text "“time” variable N". The term N_{k-1} in the numerator is annotated with a red arrow pointing to it from the text "gain of nodes of degree k". The term N_k in the numerator is annotated with a black arrow pointing to it from the text "loss of nodes of degree k". To the right of the main equation, the differential equation $\frac{dn_k}{dt} = n_{k-1} - n_k$ is shown in red.

The Degree Distribution

Master
Equation

$$\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1}$$

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“time”
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loss of nodes
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input of nodes
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The Degree Distribution

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gain of nodes of degree k

convert rate to probability

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“time” variable N
 gain of nodes of degree k
 convert rate to probability
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 input of nodes of degree 1

solving one by one:

$$N_0 = \frac{1}{N}, N_1 = \frac{N}{2}, N_2 = \frac{N}{4}, \dots$$

The Degree Distribution

Master Equation

$$\frac{dN_k}{dN} = \frac{N_{k-1} - N_k}{N} + \delta_{k,1}$$

“time” variable N → dN
gain of nodes of degree k → N_{k-1}
convert rate to probability → N
loss of nodes of degree k → N_k
input of nodes of degree 1 → $\delta_{k,1}$

$$\frac{dn_k}{dt} = n_{k-1} - n_k$$

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ansatz: $N_k \simeq N n_k$ gives

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solution: $n_k = 2^{-k}$

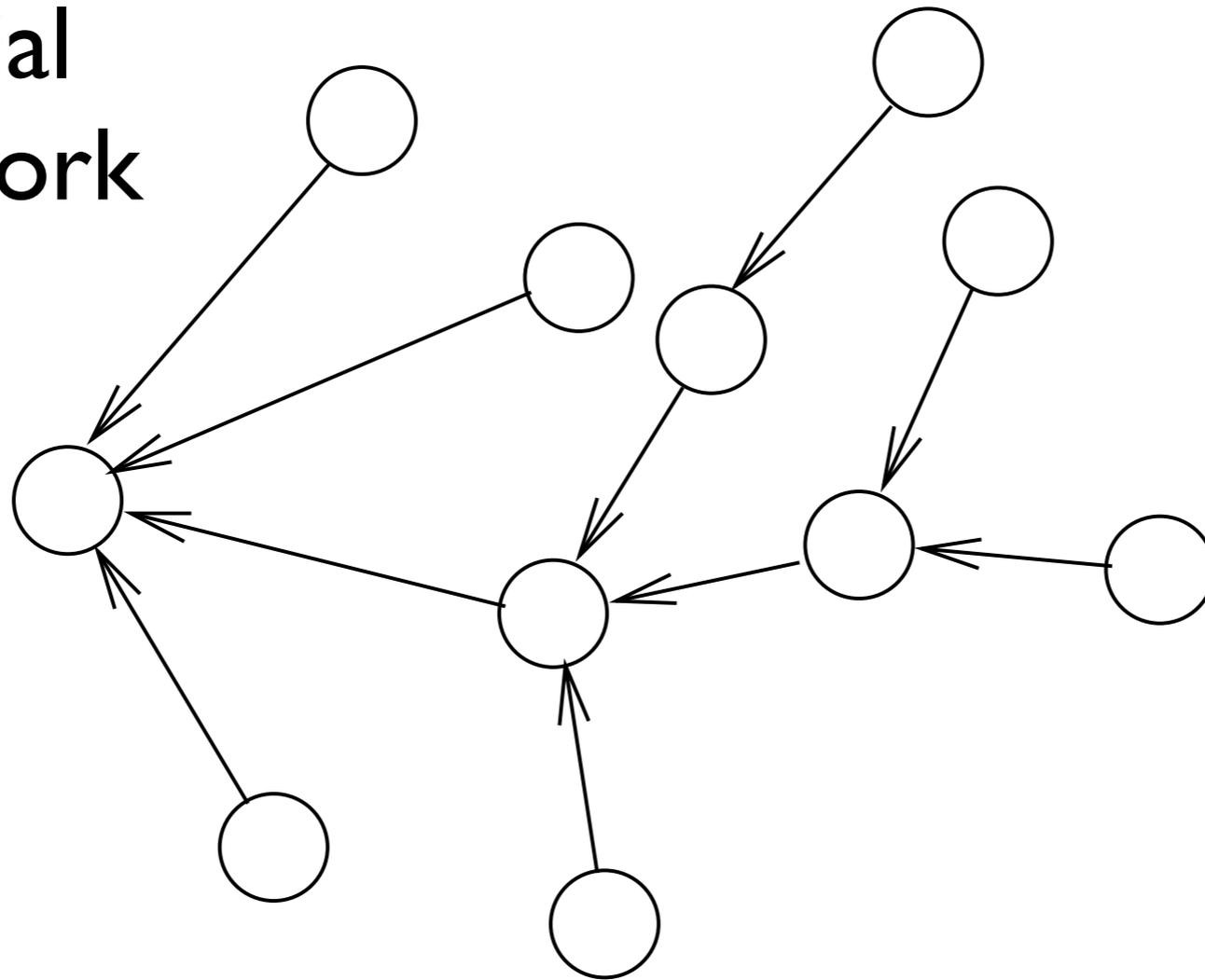
Uniform Attachment + *Redirection*

Krapivsky &
SR (2001)

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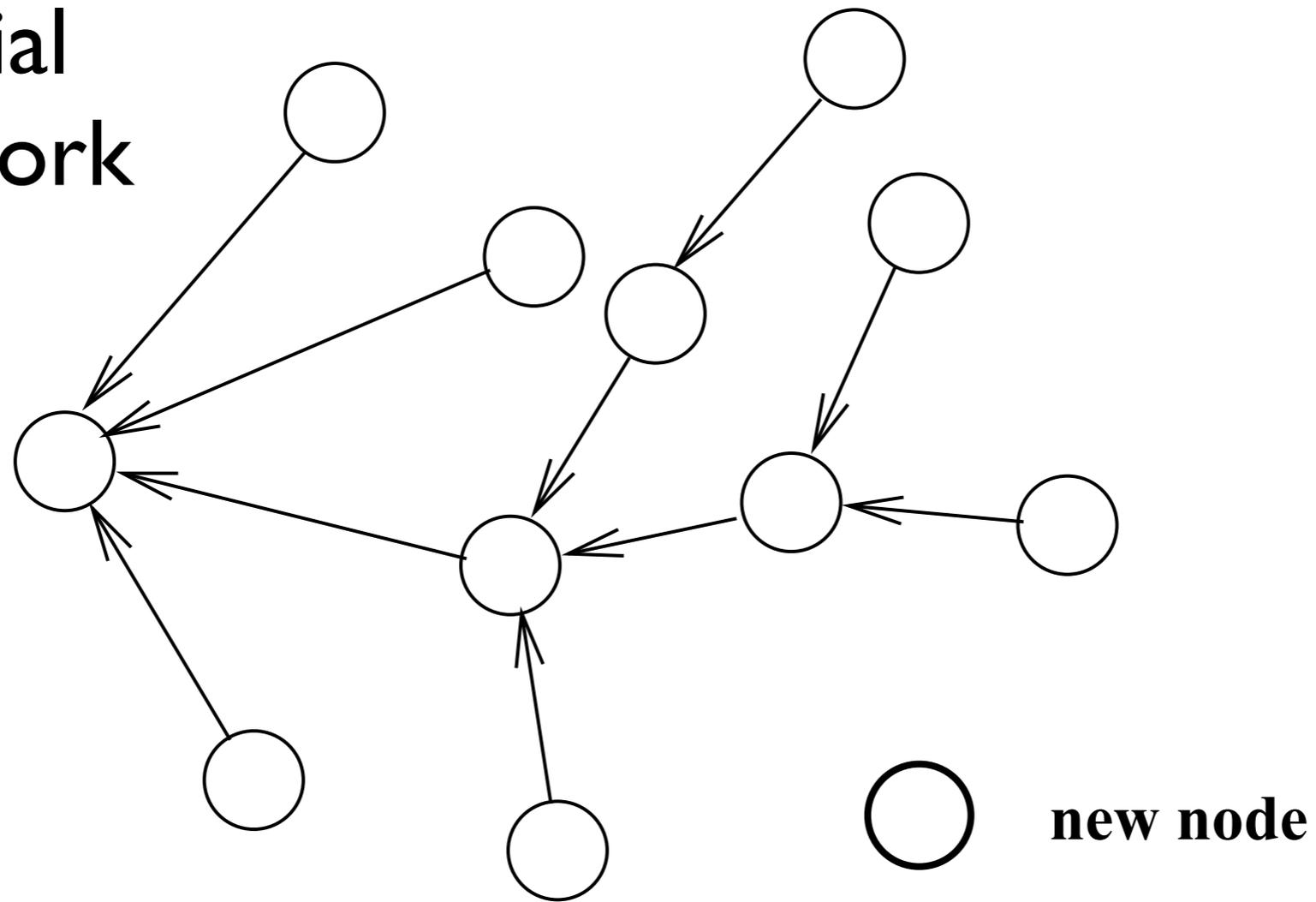
initial
network



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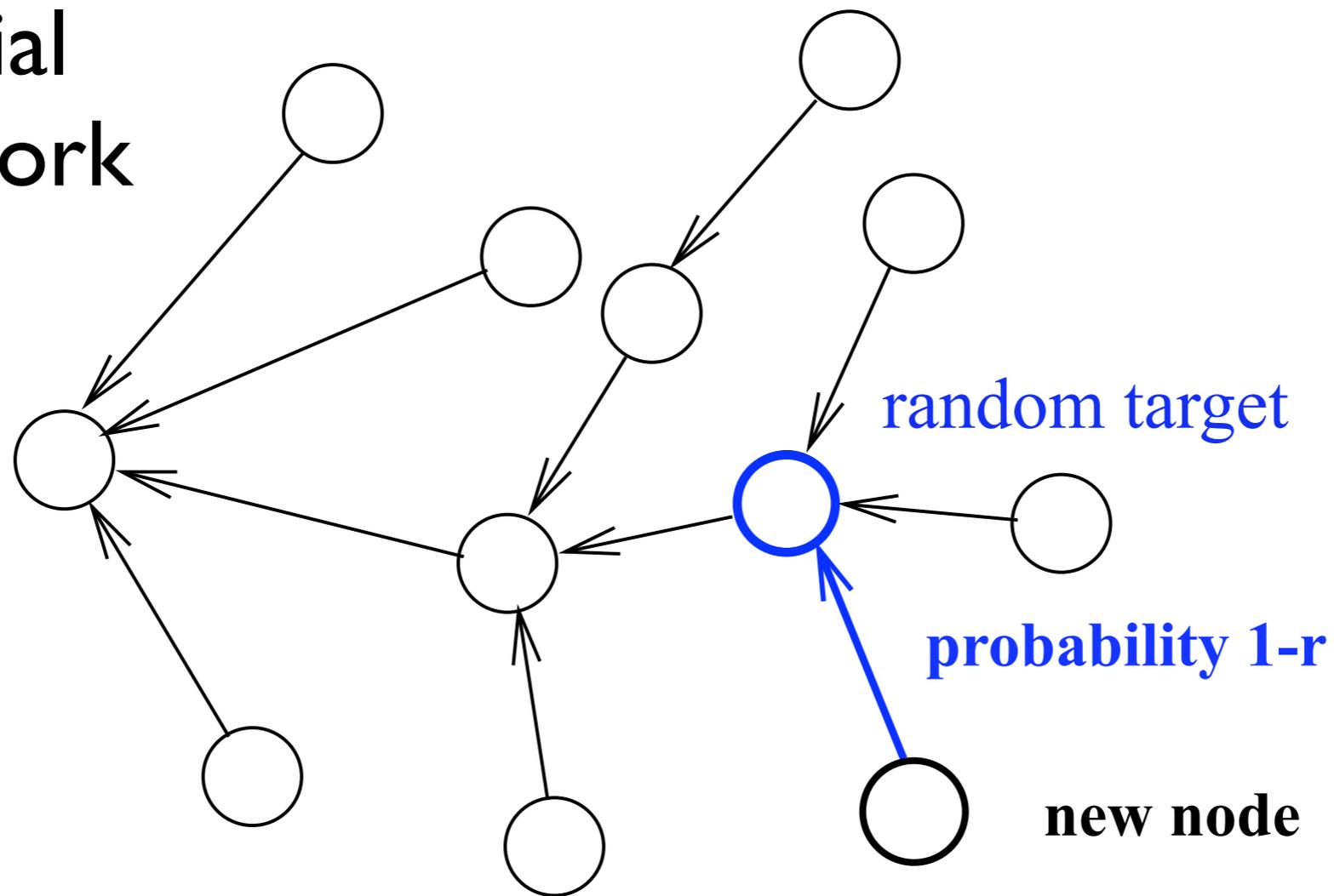
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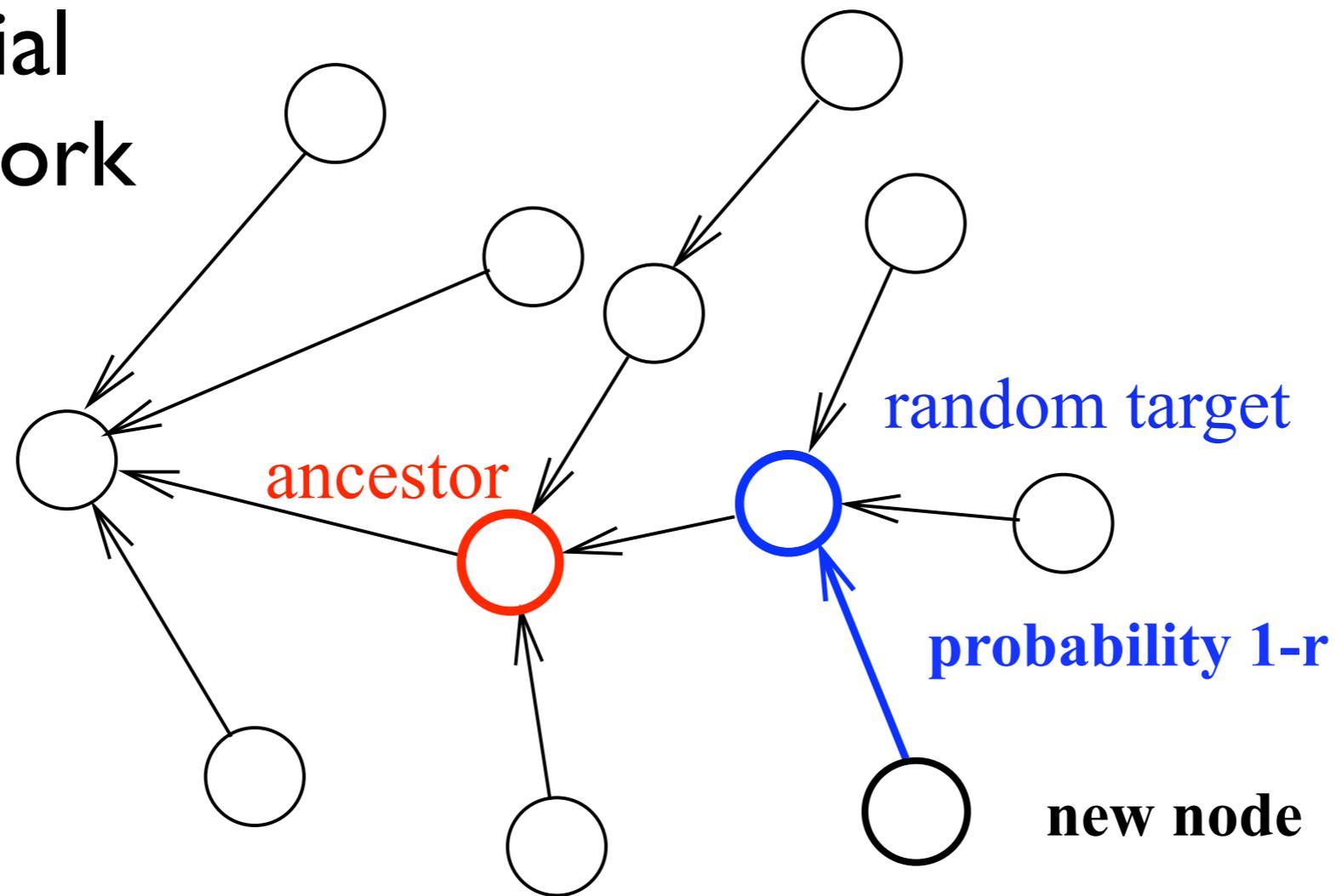
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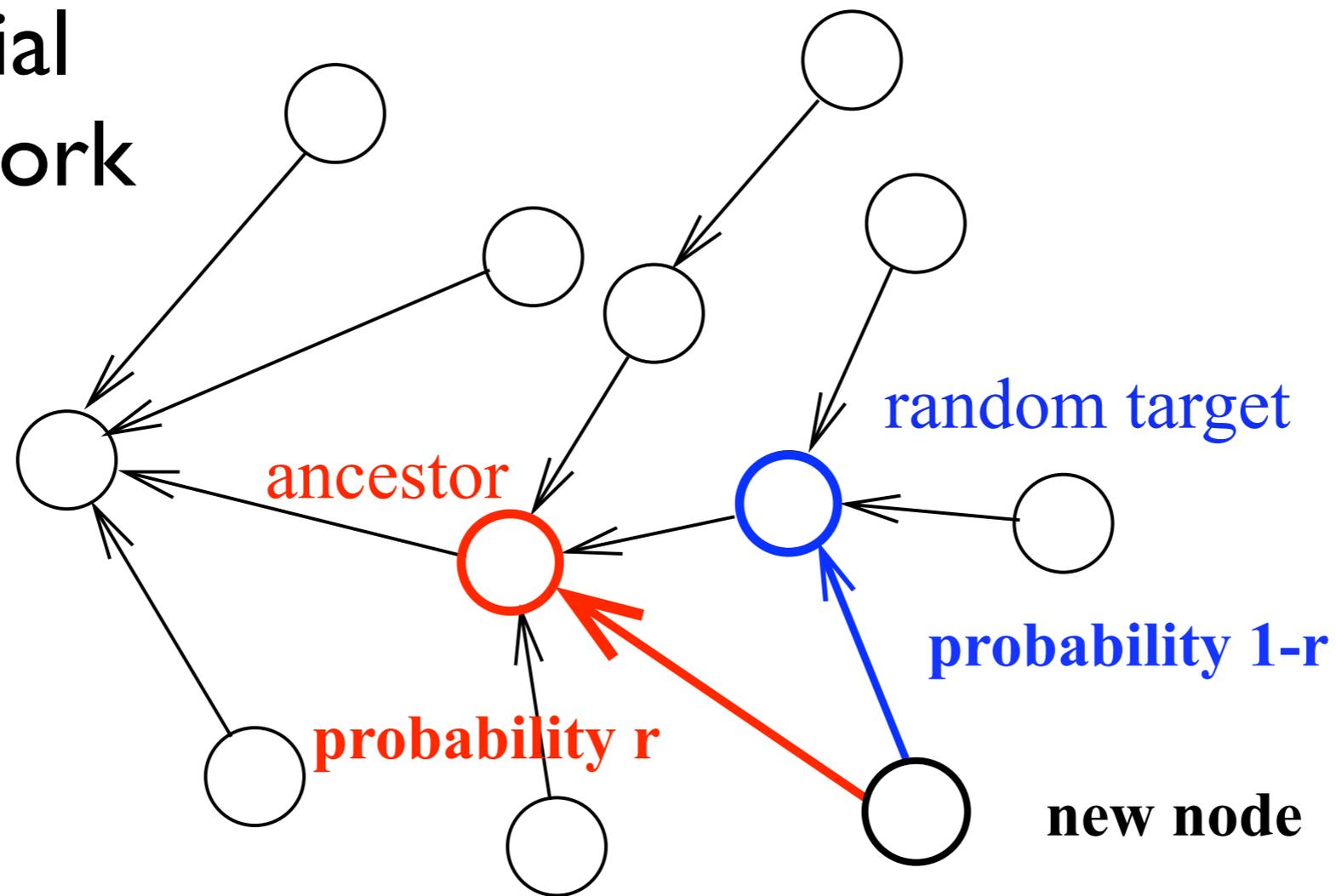
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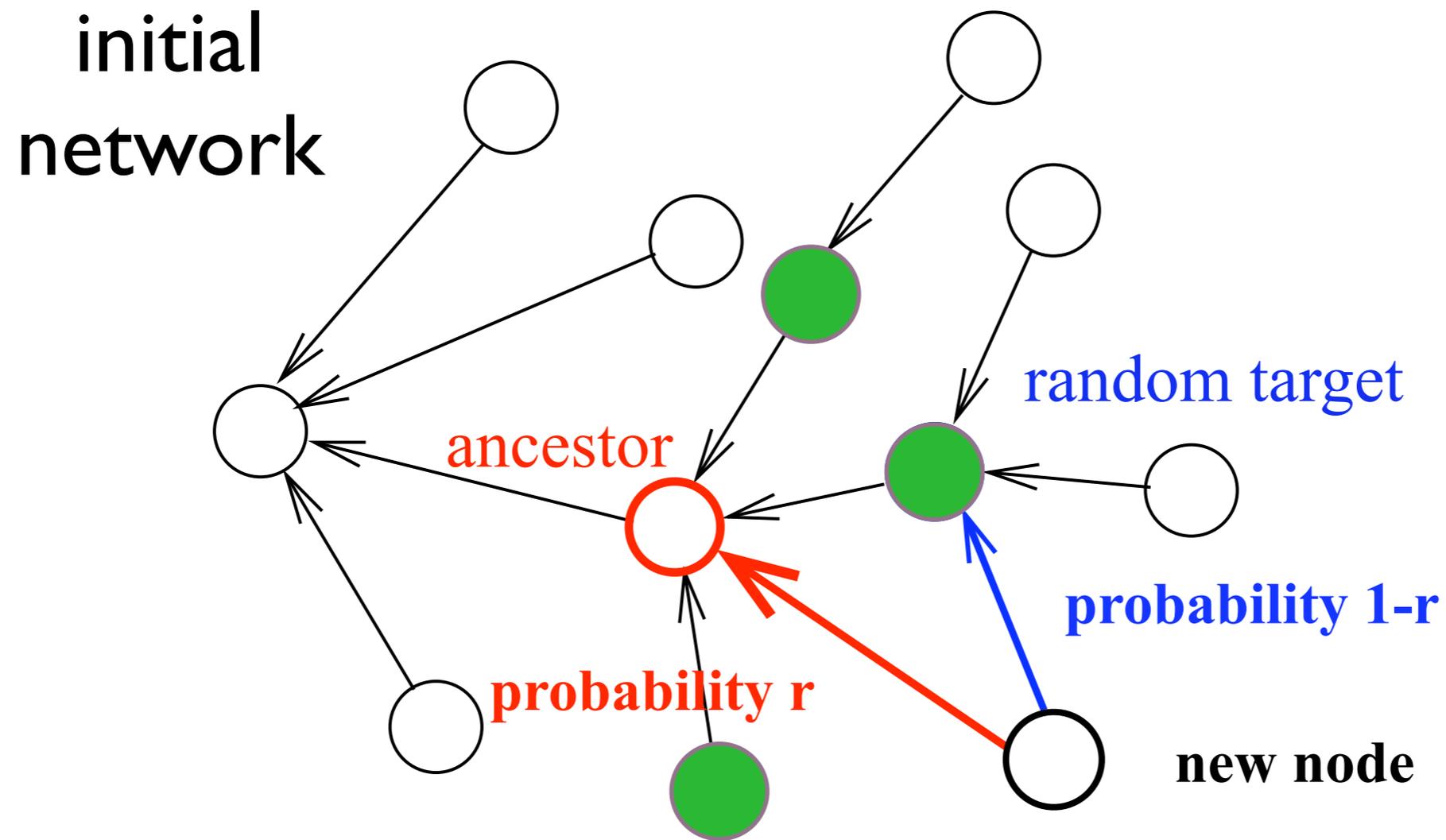
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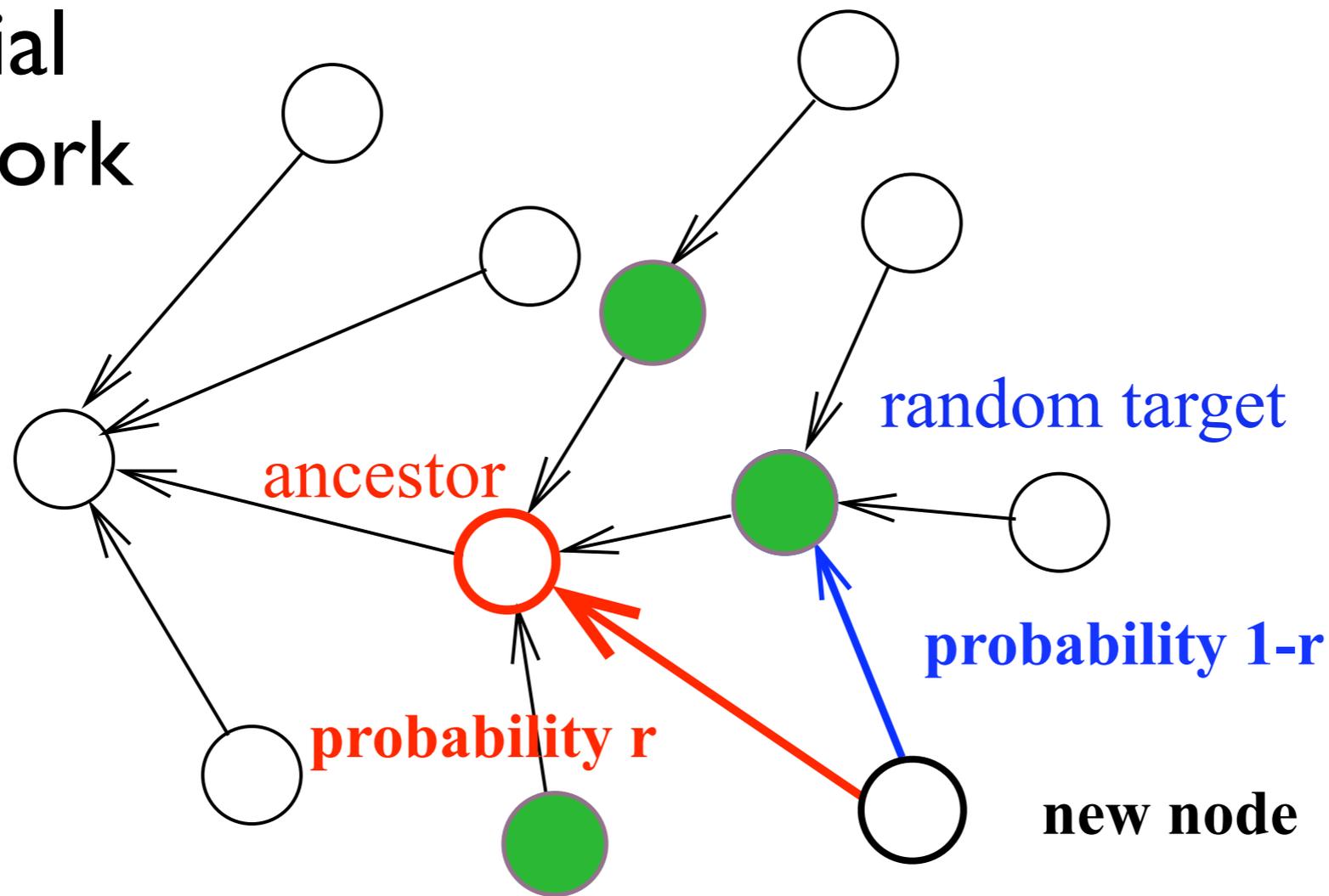
attachment rate to ancestor node:
 \propto number of upstream neighbors

Uniform Attachment + *Redirection*

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= Linear* Preferential Attachment!

initial
network



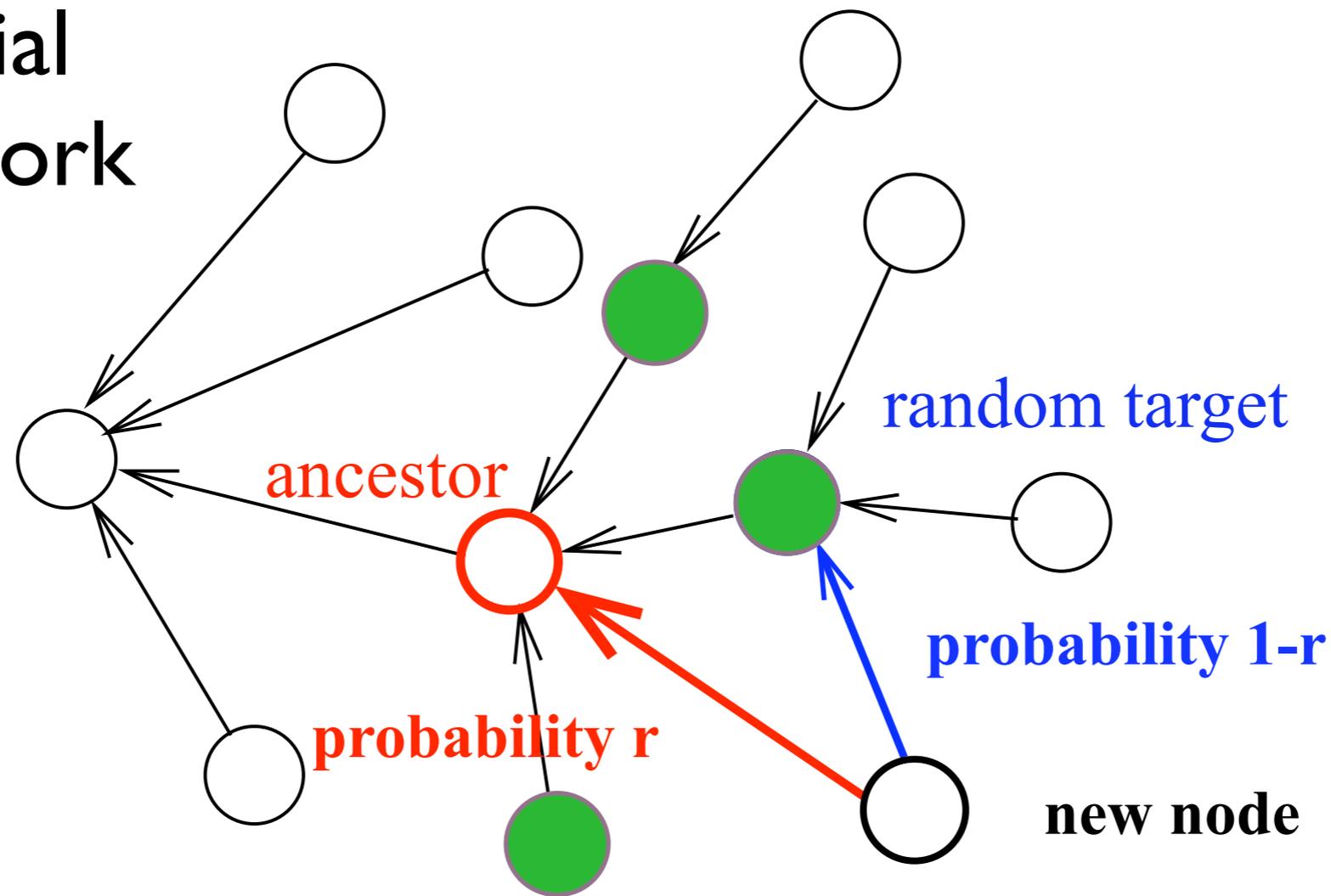
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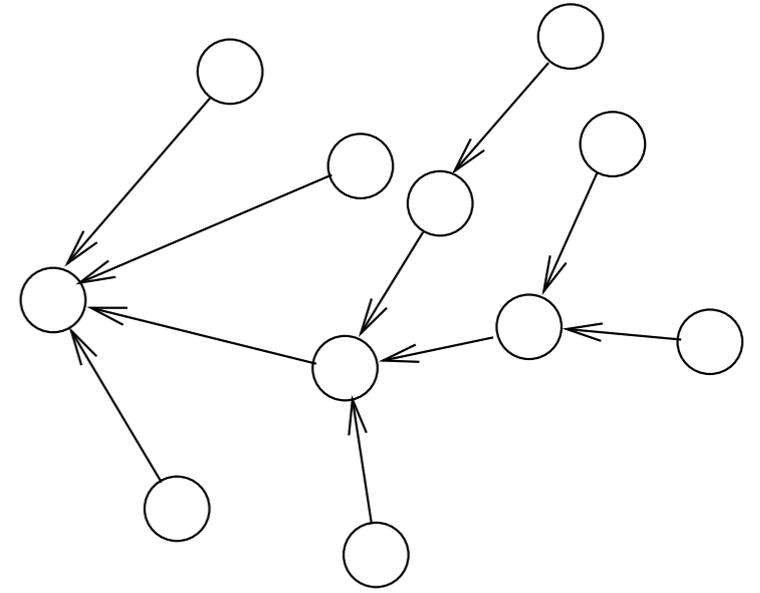
= Linear* Preferential Attachment! **shifted*

initial
network



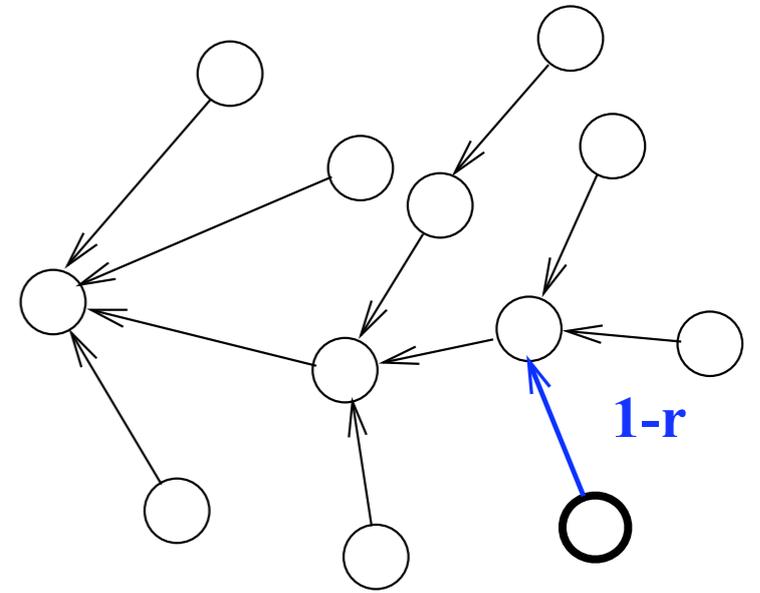
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Master Equation:



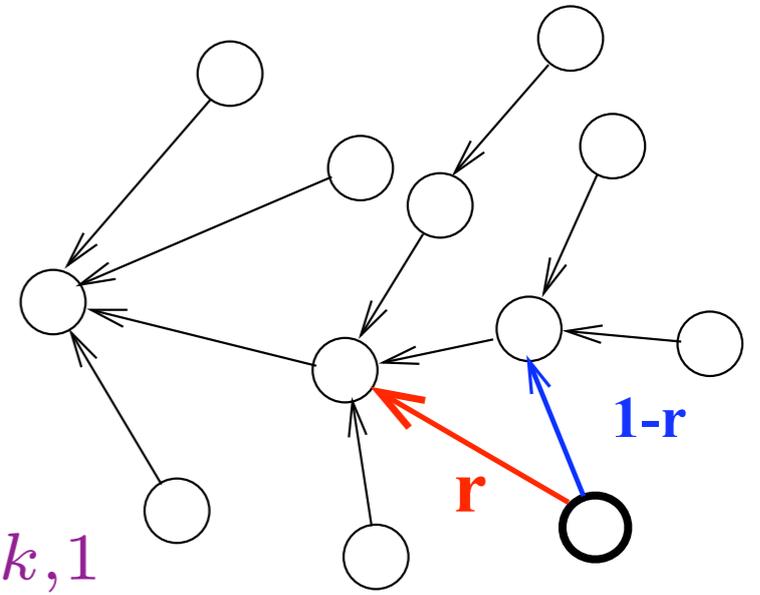
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$$+ \frac{r}{N} [(k-2)N_{k-1} - (k-1)N_k] + \delta_{k,1}$$

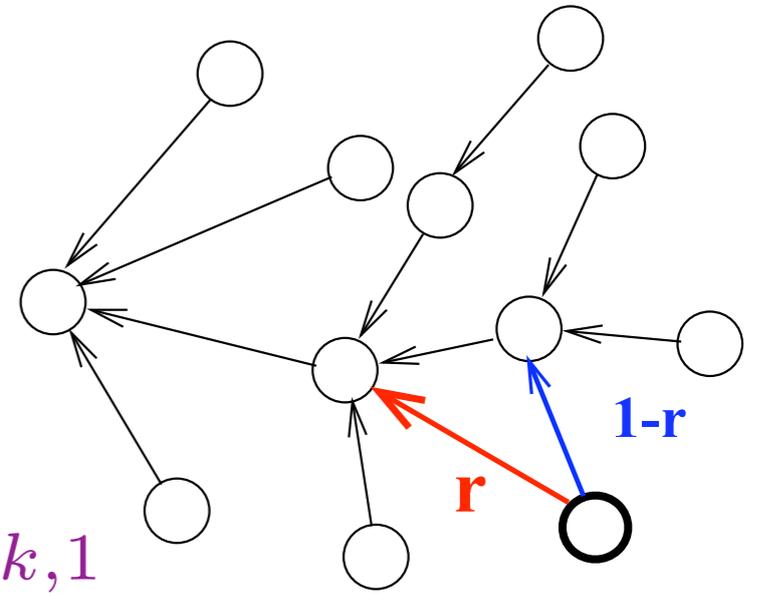


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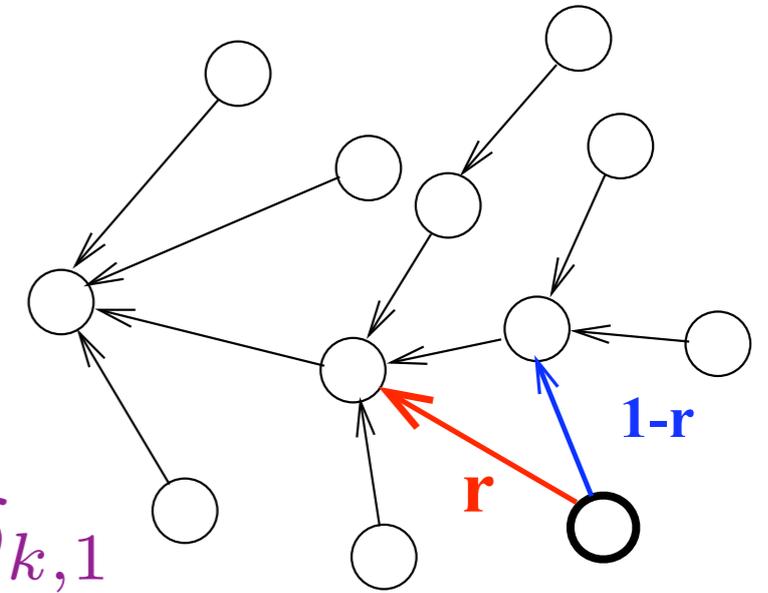
$$= \frac{r}{N} \left\{ \left[(k-1) \frac{1}{r} - 2 \right] N_{k-1} - \left[k + \frac{1}{r} + 2 \right] N_k \right\} + \delta_{k,1}$$



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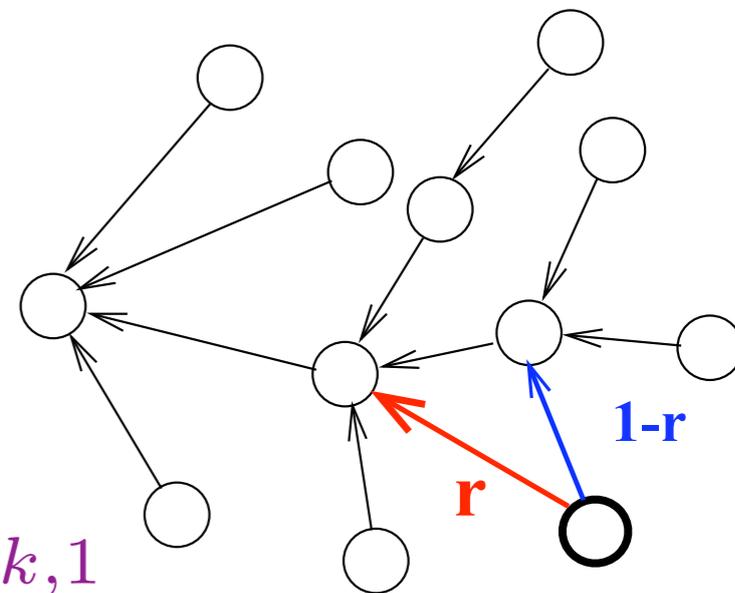
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with $A_k = k + \left(\frac{1}{r} - 2\right) \equiv k + \lambda$

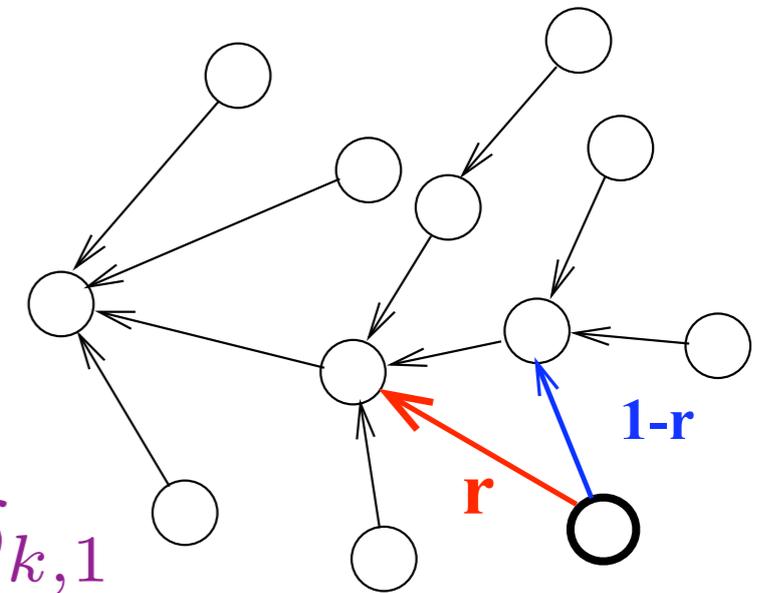
$$A = \sum_k A_k N_k = N/r$$

shifted linear attachment:

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$$A = \sum_k A_k N_k = N/r$$

shifted linear attachment:

$O(l)$ rule \rightarrow (shifted) linear preferential attachment!!

Master Equation Approach for Preferential Attachment

Krapivsky et al. (2000),
Dorogovtsev et al. (2000),
Krapivsky & SR (2001)

Master Equation:

$$\frac{dN_k}{dN} = \frac{\overset{\text{attach to node of degree } k-1}{A_{k-1} N_{k-1}}}{A} - \frac{\overset{\text{attach to node of degree } k}{A_k N_k}}{A} + \overset{\text{create node of degree } 1}{\delta_{k,1}} \quad A = \text{total rate}$$

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Degree Preferential Attachment Rate:

$$A_k \sim k^\gamma$$

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Degree Preferential Attachment Rate:

$$A_k \sim k^\gamma$$

Total Rate:

$$A = A(N) = \sum_{j=1}^{\infty} A_j N_j = \sum_{j=1}^{\infty} j^\gamma N_j \equiv M_\gamma(N)$$

Moment equations:

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These suggest: $A(N) = \sum_j j^\gamma N_j \propto \mu(\gamma) N$ for $0 \leq \gamma \leq 1$

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Converts rate eqns. to linear recursions

$$\frac{dN_k}{dN} = \frac{A_{k-1}N_{k-1}}{A} - \frac{A_k N_k}{A} + \delta_{k,1}$$

$$\Rightarrow n_k = \frac{A_{k-1}n_{k-1} - A_k n_k}{\mu} + \delta_{k,1}$$

Formal Solution

$$n_k = \frac{A_{k-1}n_{k-1} - A_k n_k}{\mu} + \delta_{k,1}$$

Formal Solution

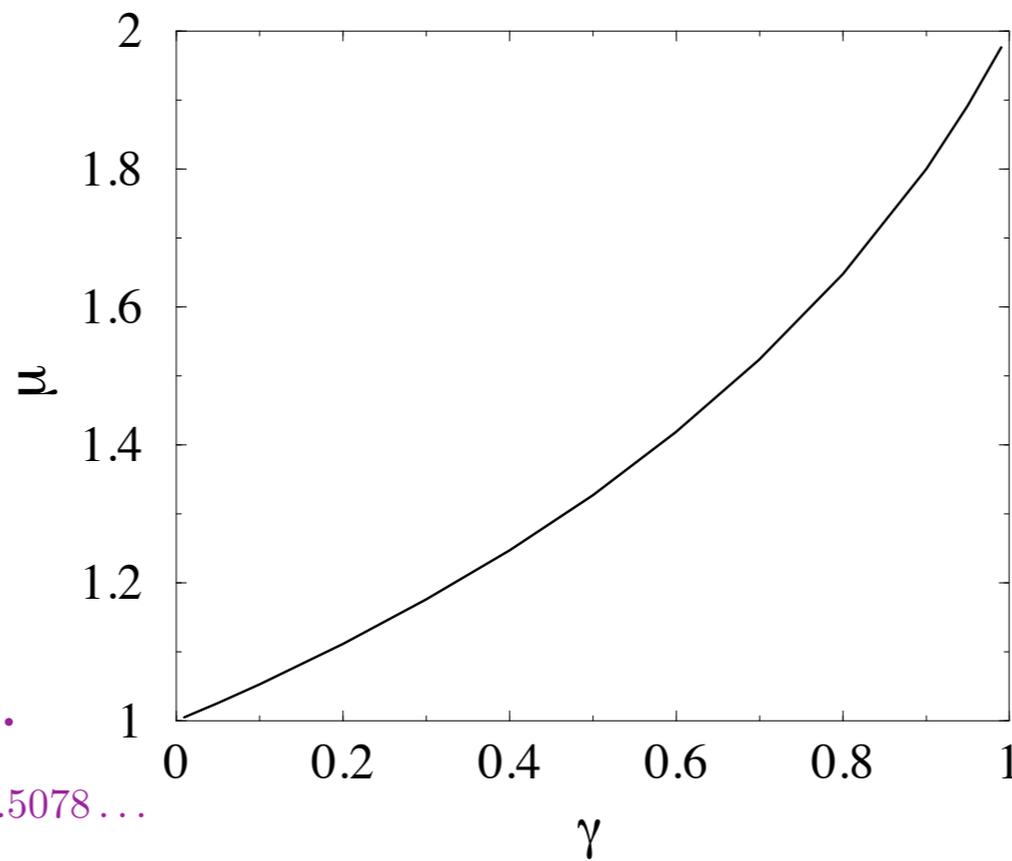
$$n_k = \frac{A_{k-1}n_{k-1} - A_k n_k}{\mu} + \delta_{k,1} \longrightarrow n_k = \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j} \right)^{-1}$$

Formal Solution

$$n_k = \frac{A_{k-1}n_{k-1} - A_k n_k}{\mu} + \delta_{k,1} \longrightarrow n_k = \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j}\right)^{-1}$$

determination of $\mu = \sum_{k=1}^{\infty} A_k n_k$: $\longrightarrow 1 = \sum_{k=1}^{\infty} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j}\right)^{-1}$

numerical solution for $A_k = k^\gamma$



$$\mu = 2 - C(1 - \gamma) + \dots$$

$$C = 4 \sum_{j=1}^{\infty} \frac{\ln j}{(j+1)(j+2)} = 2.407\dots$$

$$\mu = 1 + B\gamma + \dots$$

$$B = \sum_{j=1}^{\infty} \frac{\ln j}{2^j} = 0.5078\dots$$

Formal Solution: $n_k = \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j} \right)^{-1}$

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Asymptotics for $A_k \sim k^\gamma$

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Asymptotics for $A_k \sim k^\gamma$ (use $\prod_j (1 + \epsilon_j)^{-1} \sim e^{-\int^k \epsilon_j dj}$)

$$n_k \sim \left\{ \begin{array}{l} k^{-\gamma} e^{-[\mu k^{1-\gamma}/(1-\gamma)]} \quad 0 \leq \gamma < 1 \end{array} \right.$$

Formal Solution: $n_k = \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j} \right)^{-1}$

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$$k^{-\gamma} e^{-[\mu k^{1-\gamma}/(1-\gamma)]} \quad 0 \leq \gamma < 1$$

universal, generic

$n_k \sim$

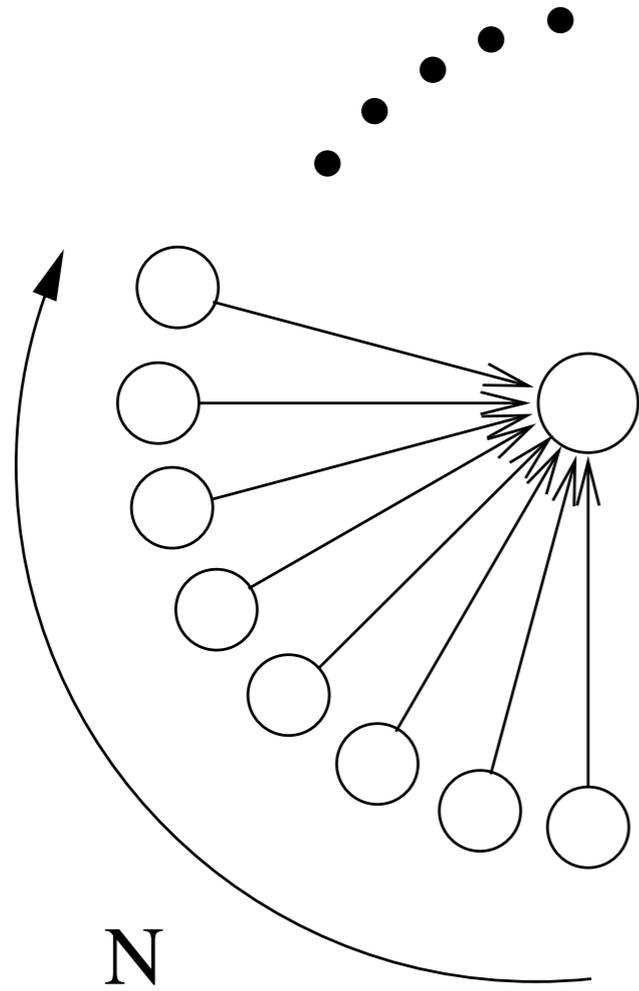


Formal Solution: $n_k = \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j} \right)^{-1}$

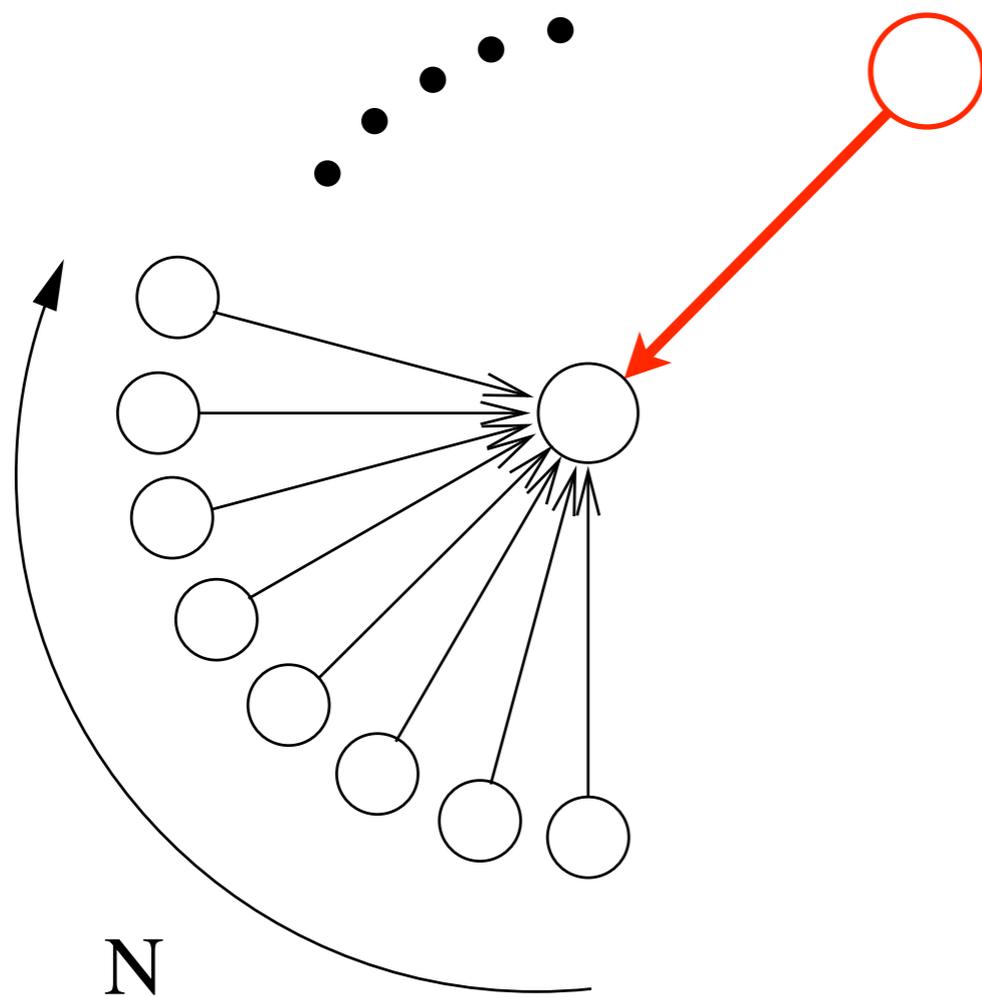
Asymptotics for $A_k \sim k^\gamma$

$$n_k \sim \begin{cases} k^{-\gamma} e^{-[\mu k^{1-\gamma}/(1-\gamma)]} & 0 \leq \gamma < 1 \\ \text{"bible"} & \gamma > 2 \end{cases}$$

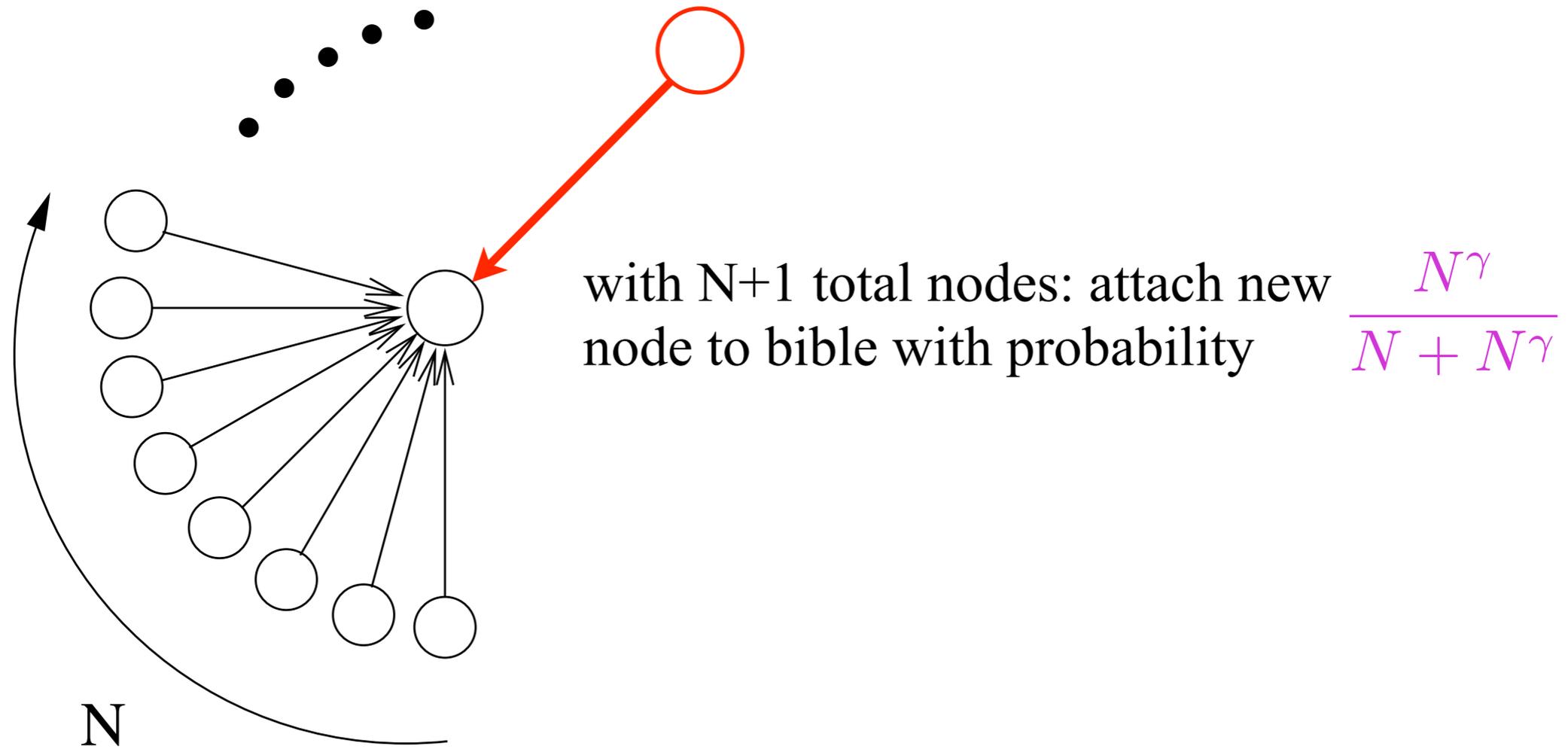
Creating a “Bible”



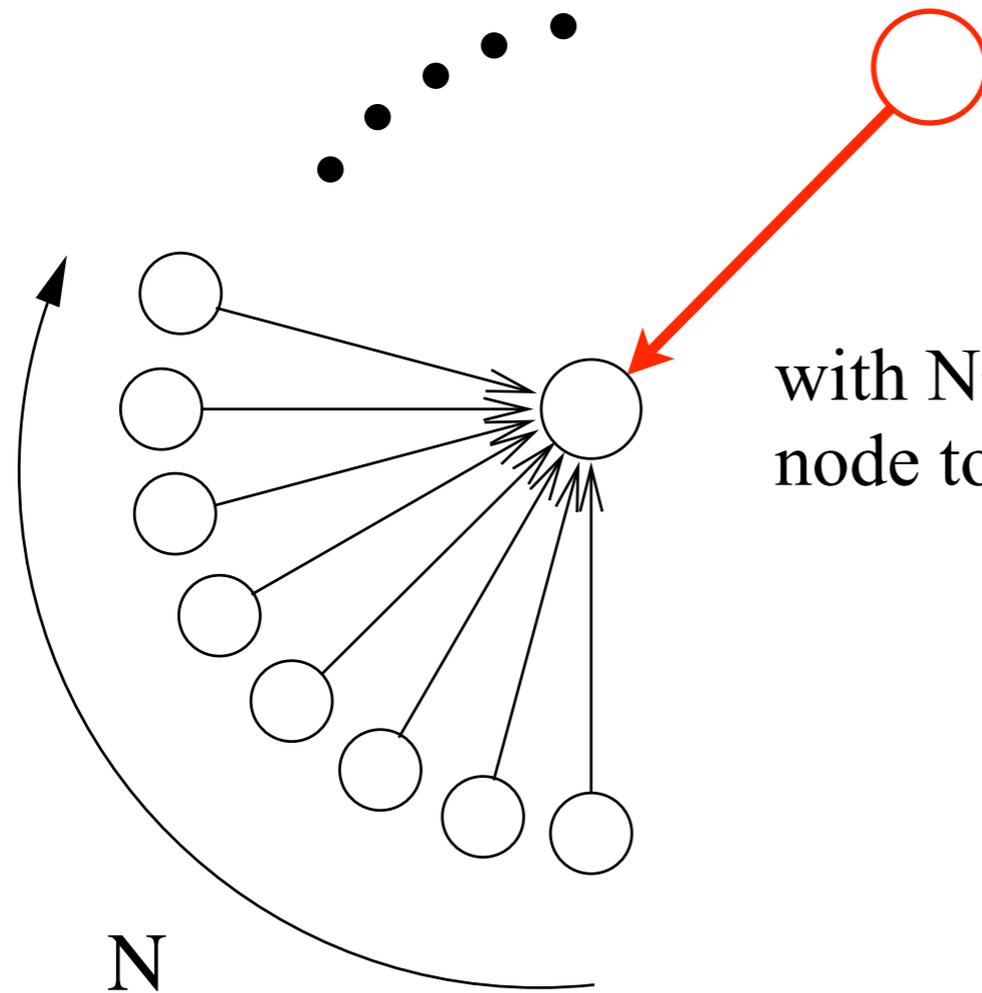
Creating a “Bible”



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with $N+1$ total nodes: attach new node to bible with probability

$$\frac{N^\gamma}{N + N^\gamma}$$

bible probability: $\mathcal{P} = \prod_{N=1}^{\infty} \frac{1}{1 + N^{1-\gamma}} = \begin{cases} \text{zero} & \gamma \leq 2 \\ \text{non-zero} & \gamma > 2 \end{cases}$

Formal Solution: $n_k = \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j} \right)^{-1}$

Asymptotics for $A_k \sim k^\gamma$

$$n_k \sim \begin{cases} k^{-\gamma} e^{-[\mu k^{1-\gamma}/(1-\gamma)]} & 0 \leq \gamma < 1 \\ \text{"bible"} & \gamma > 2 \end{cases}$$

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Origin of Non-Universal Degree Distributions

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$$\longrightarrow 1 < \mu < \infty$$

$$2 < \nu < \infty$$

Linear Attachment $A_k = k$:

$$n_k = \frac{4}{k(k+1)(k+2)}$$

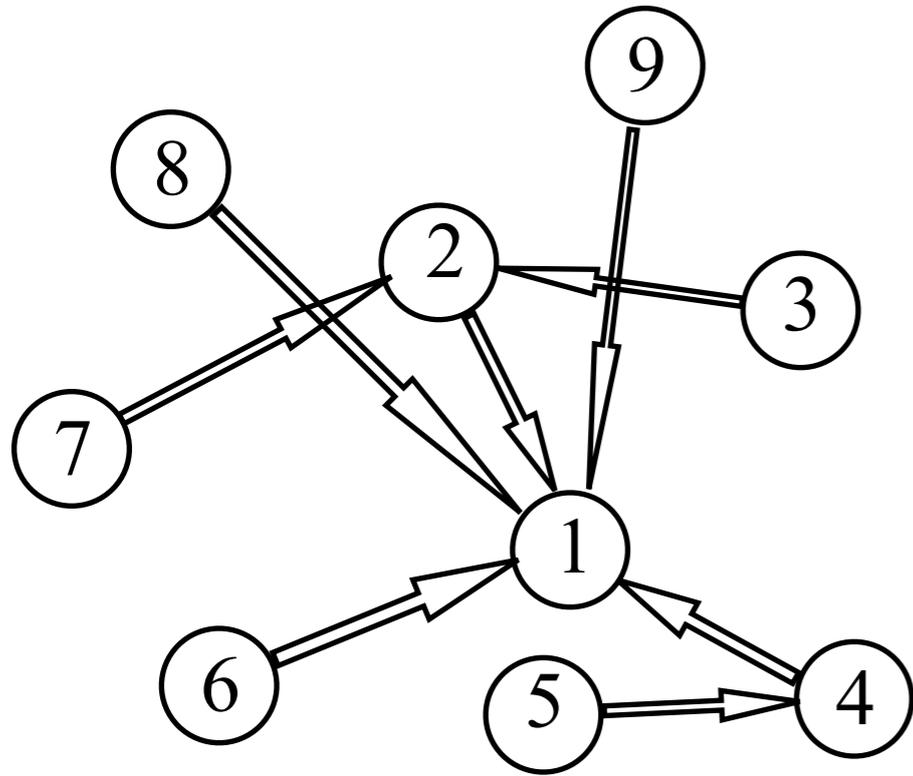
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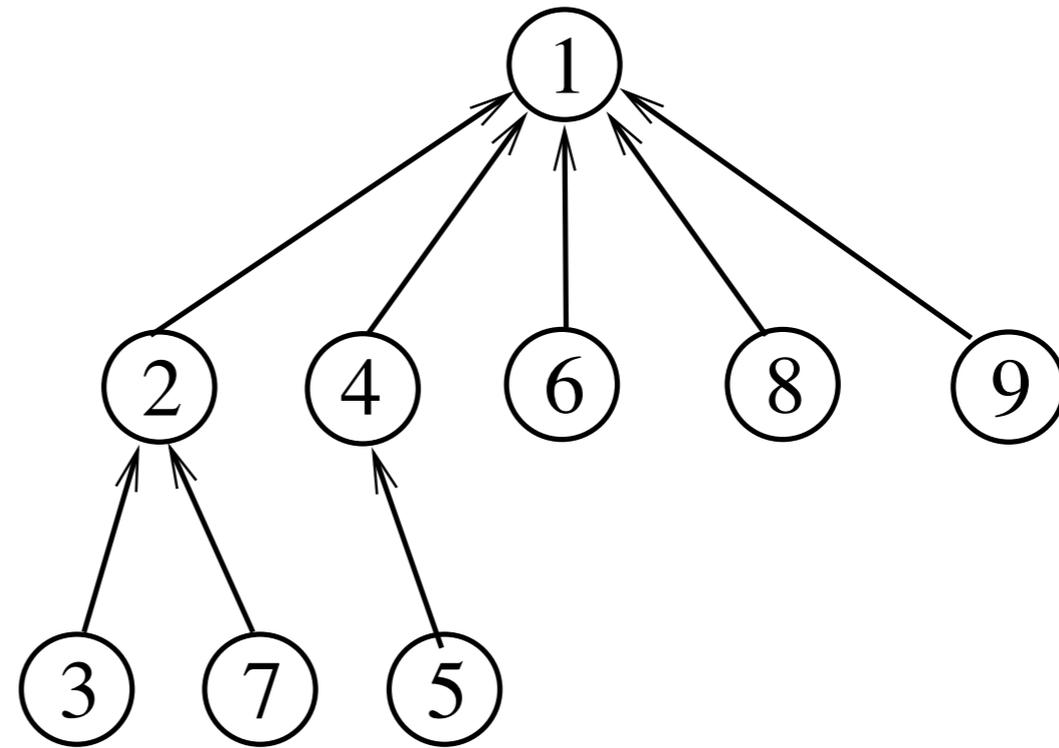
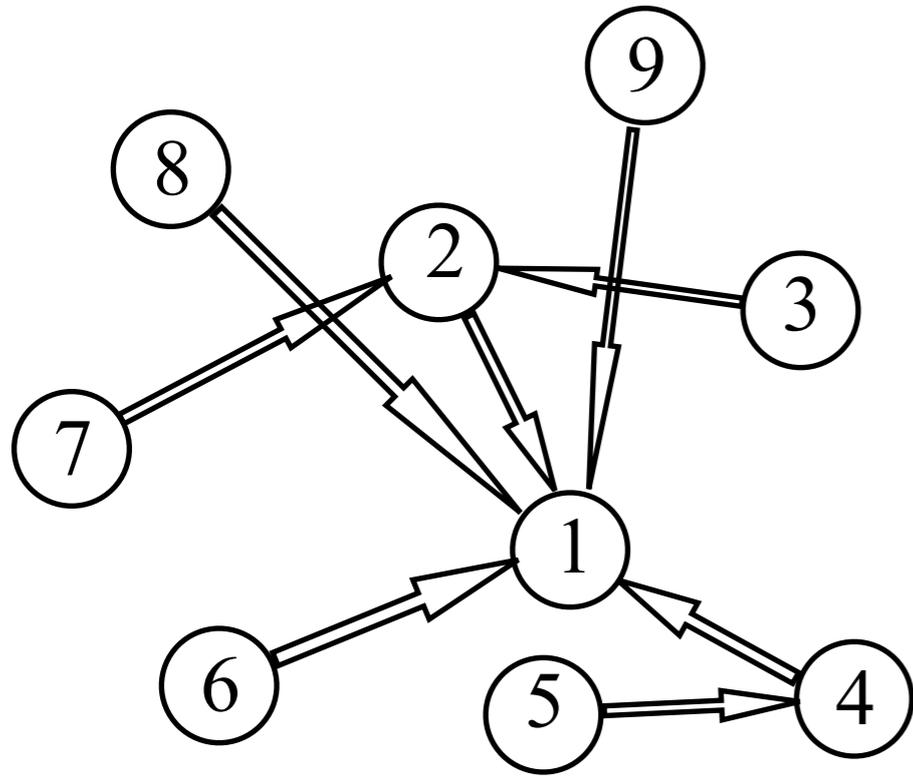
Shifted Linear Attachment $A_k = k + \lambda$

$$\begin{aligned} n_k &= \frac{\mu}{A_k} \prod_{j=1}^k \left(1 + \frac{\mu}{A_j} \right)^{-1} \\ &= (2 + \lambda) \frac{\Gamma(3 + 2\lambda)}{\Gamma(1 + \lambda)} \frac{\Gamma(k + \lambda)}{\Gamma(k + 3 + 2\lambda)} \\ &\sim k^{-(3+\lambda)} \quad -1 < \lambda < \infty \end{aligned}$$

Genealogy (for uniform attachment)



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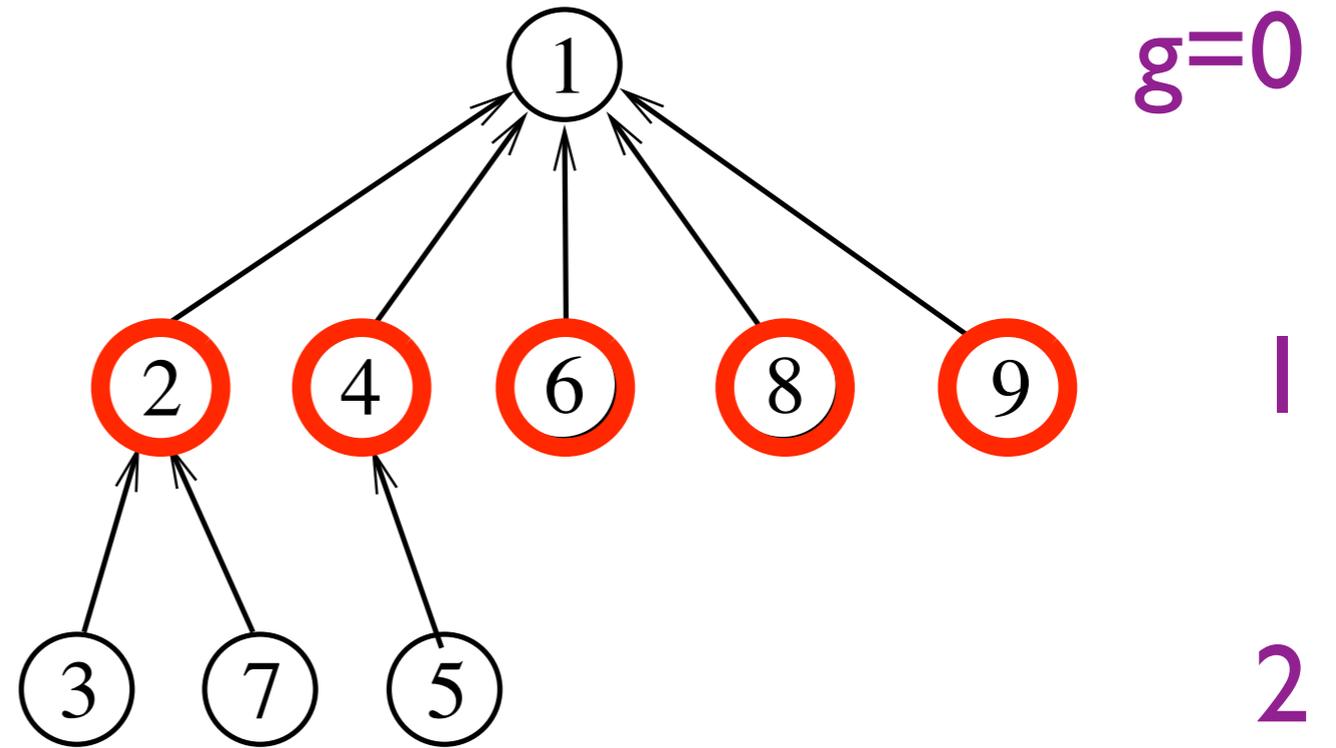
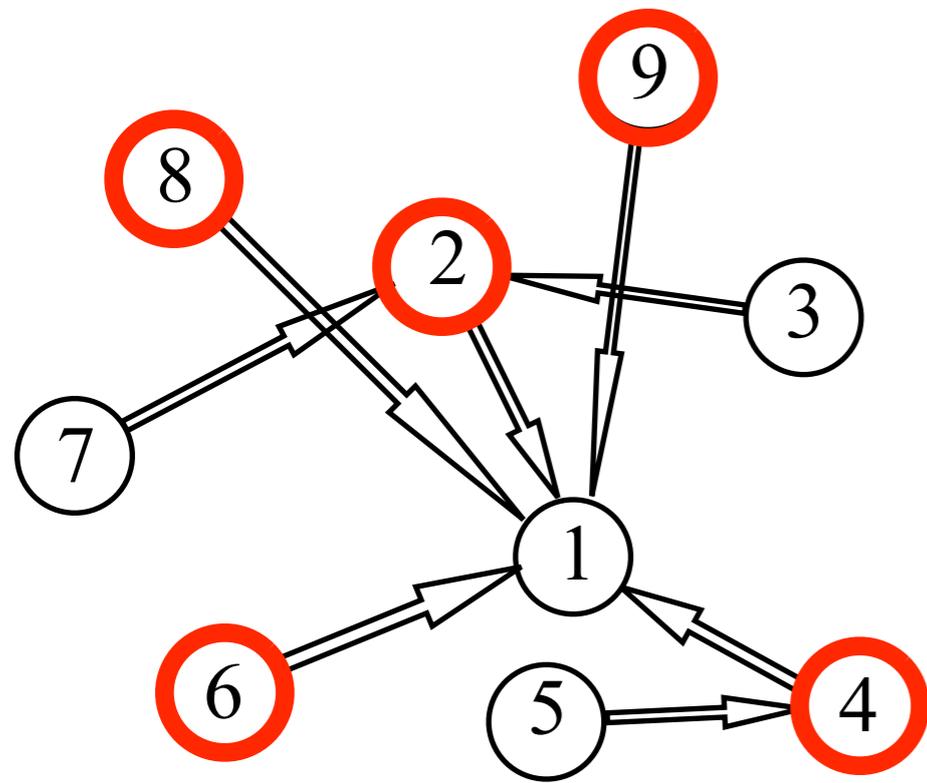


$g=0$

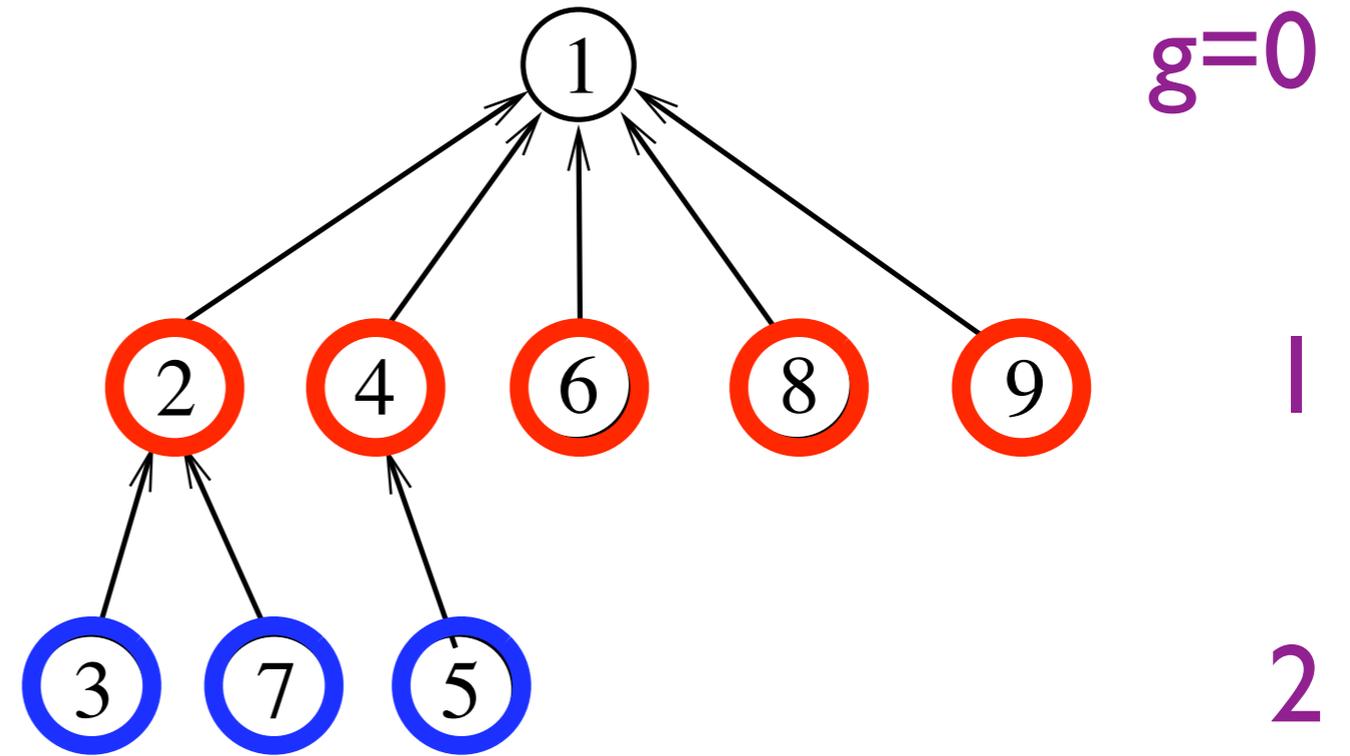
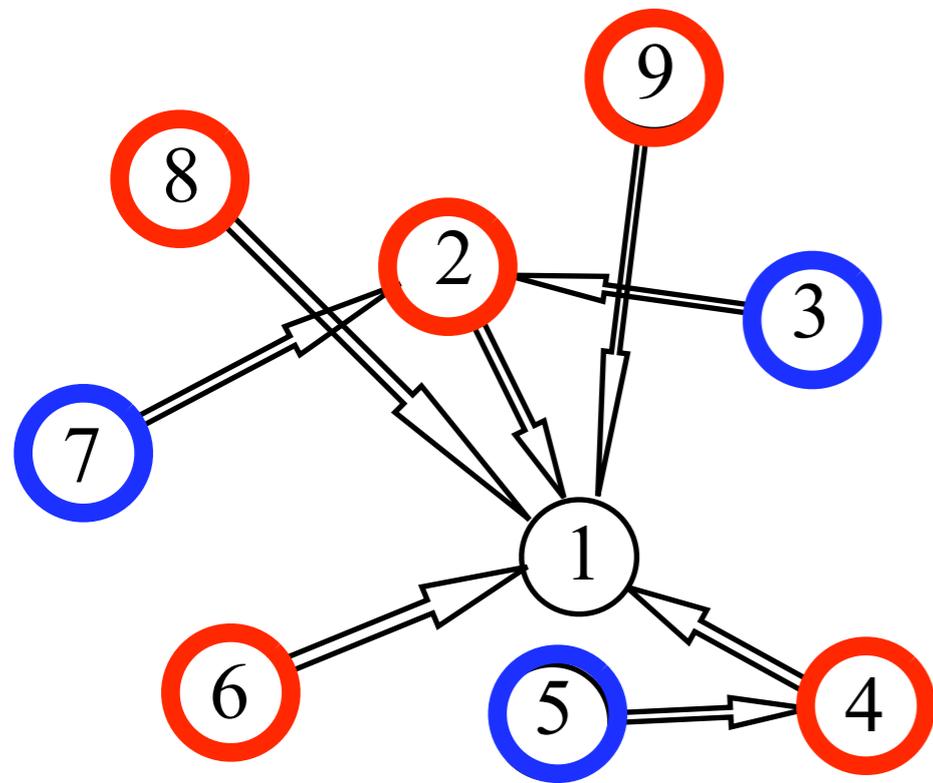
1

2

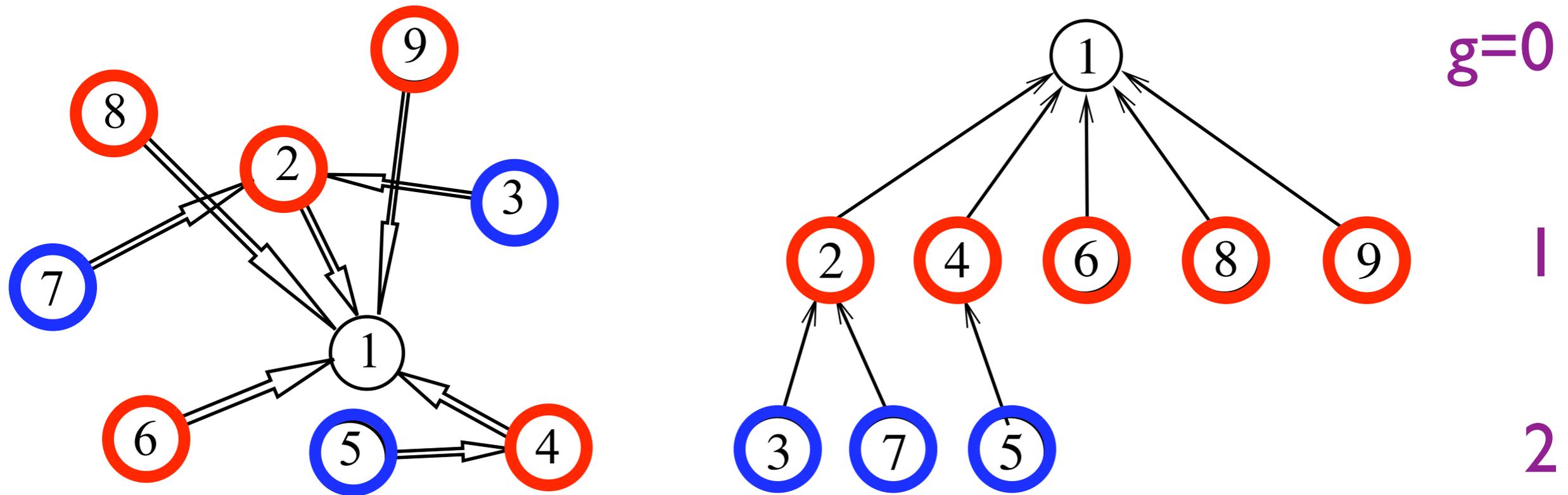
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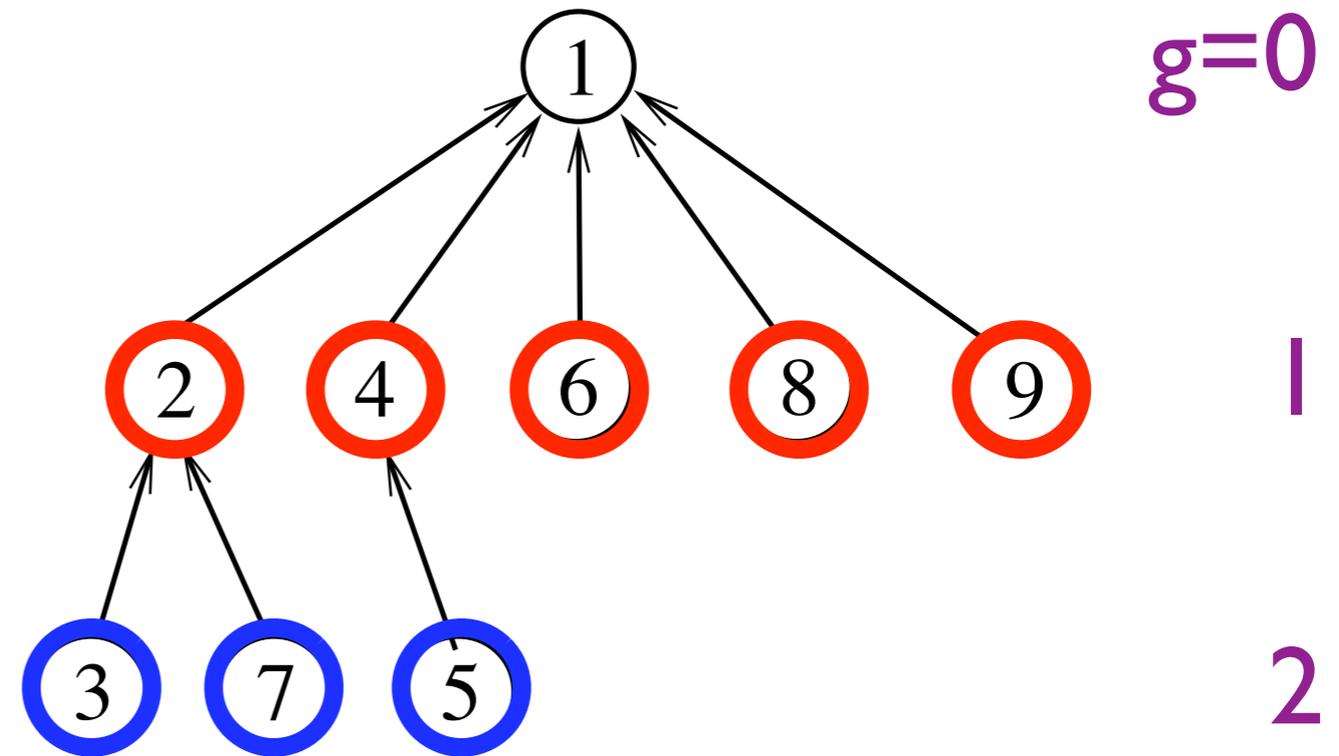
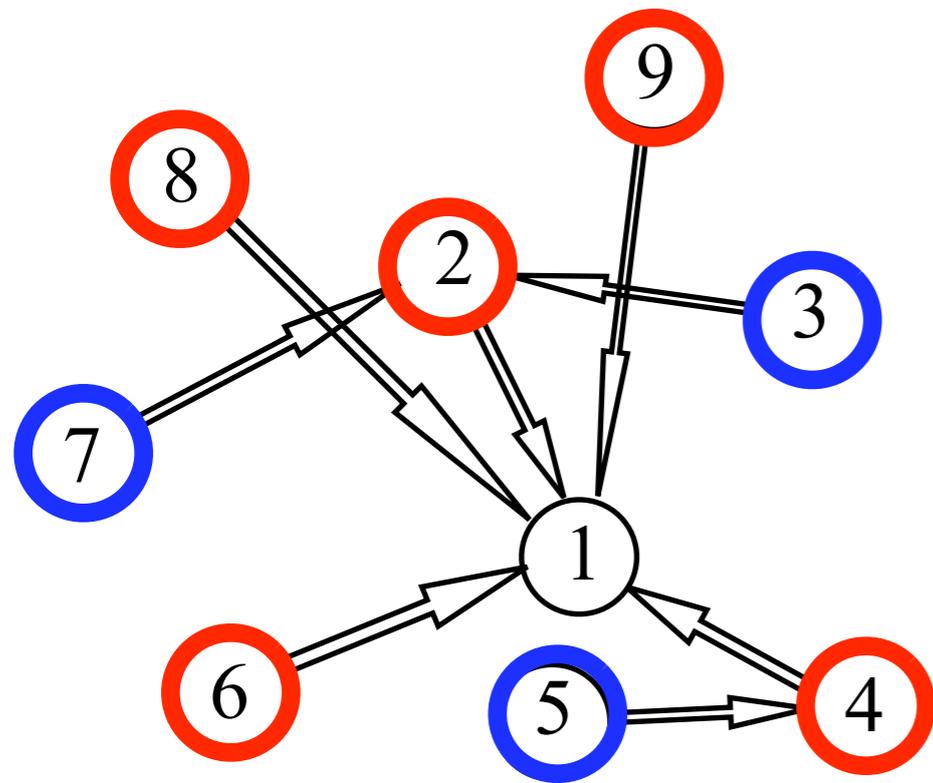


Genealogy (for uniform attachment)



Size of generation $g \equiv L_g$

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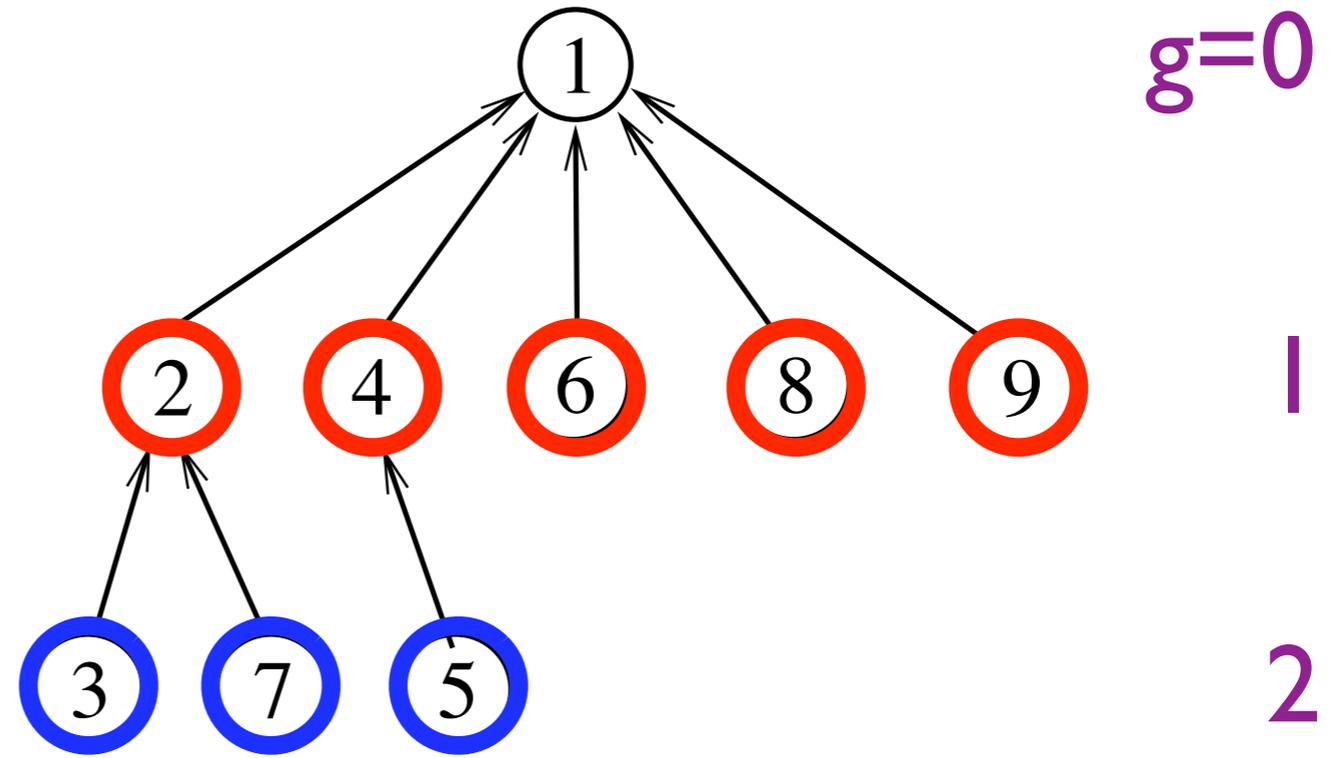
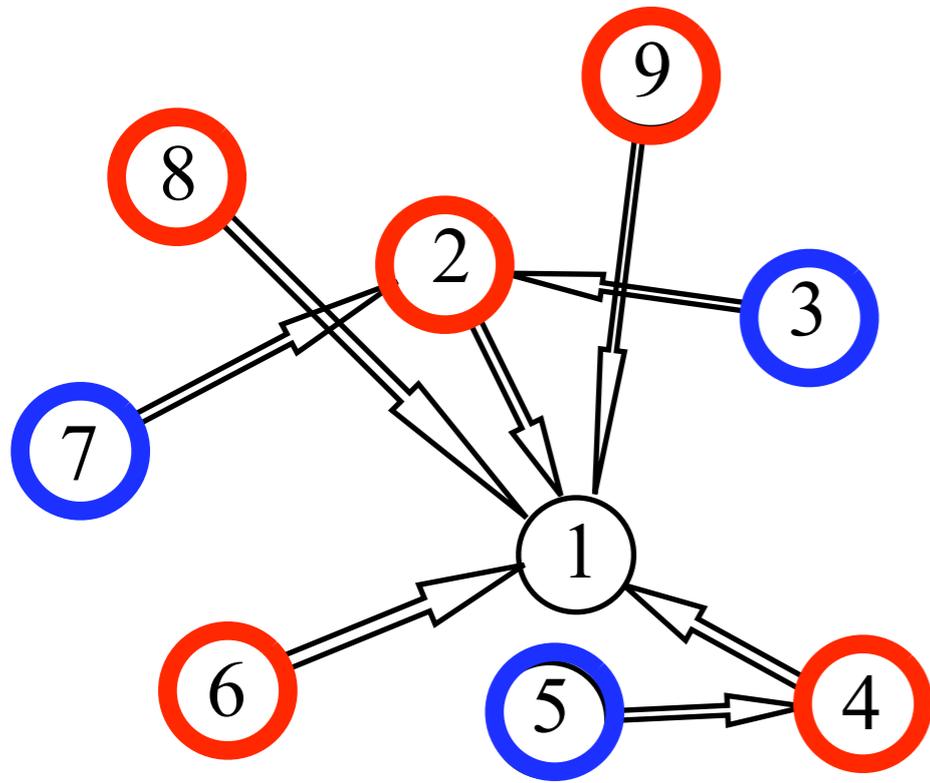


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rate equation:

$$\frac{dL_g}{dN} = \frac{L_{g-1}}{N}$$

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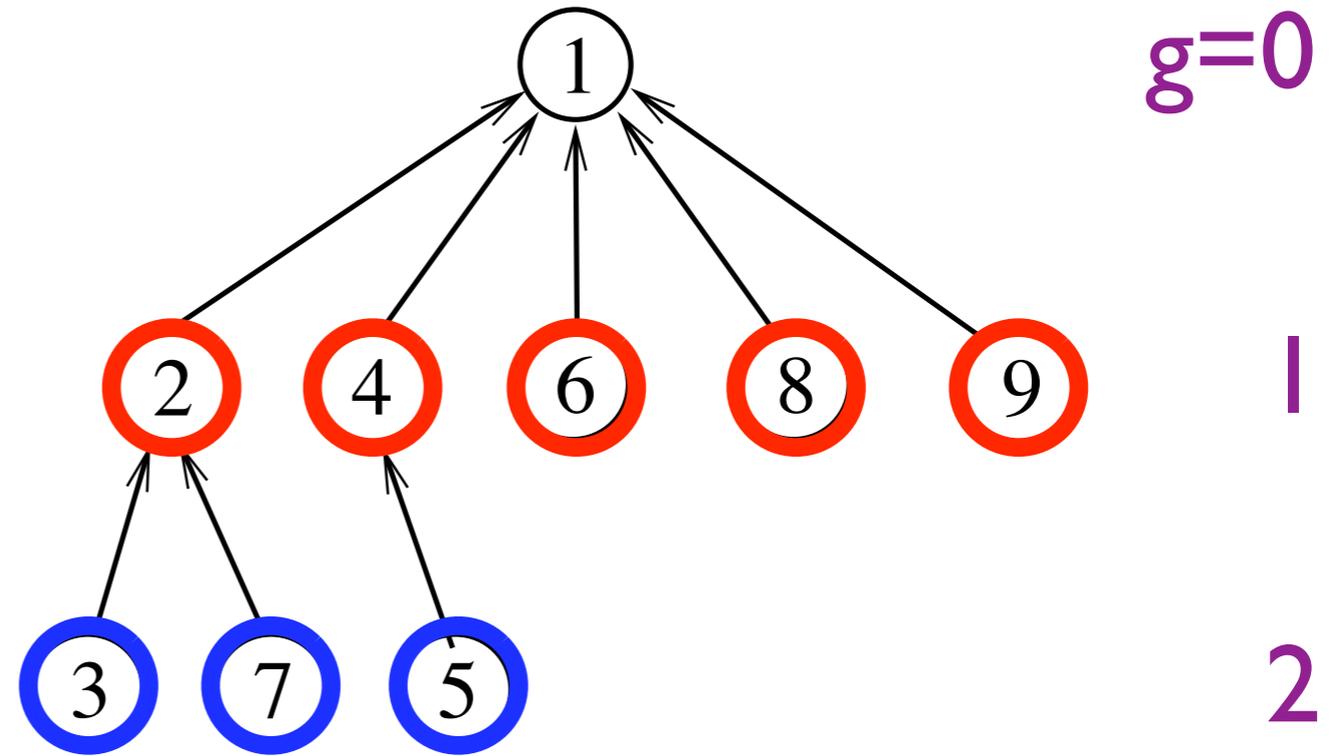
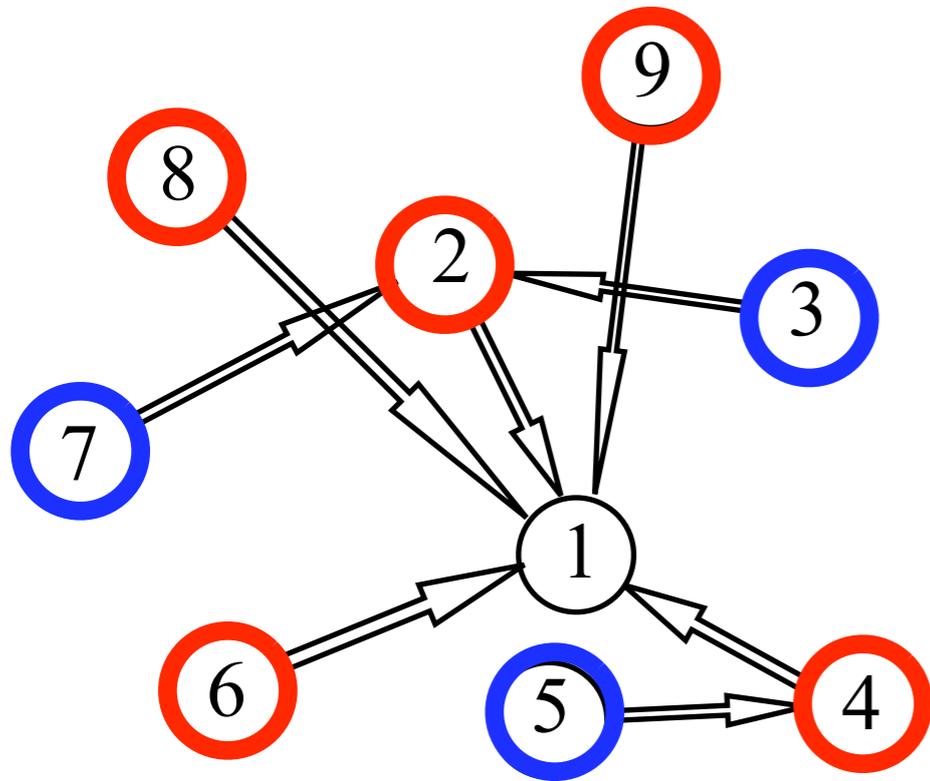
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Solution: $L_g = \frac{(\ln N)^g}{g!}$

$L_g = 1$ defines last generation

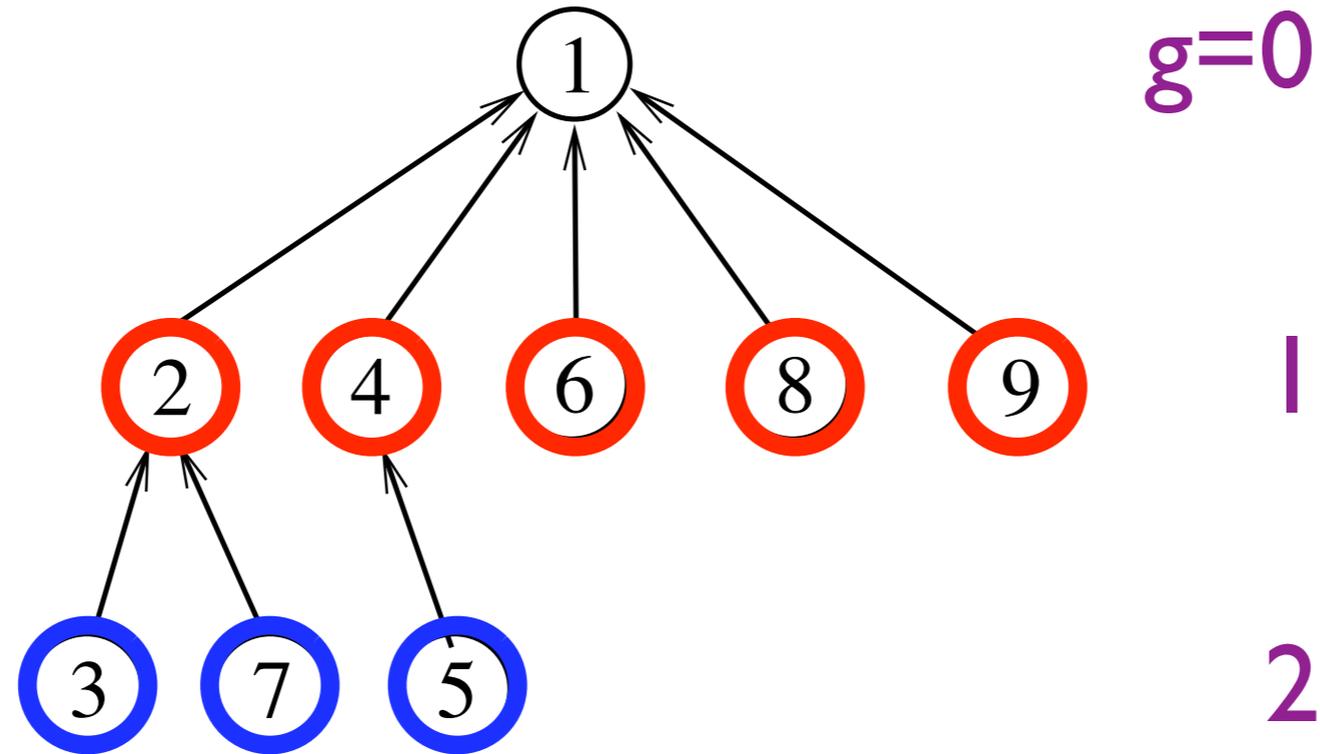
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$\longrightarrow g_{\max} \sim e \ln N$

diameter $\sim 2e \ln N$

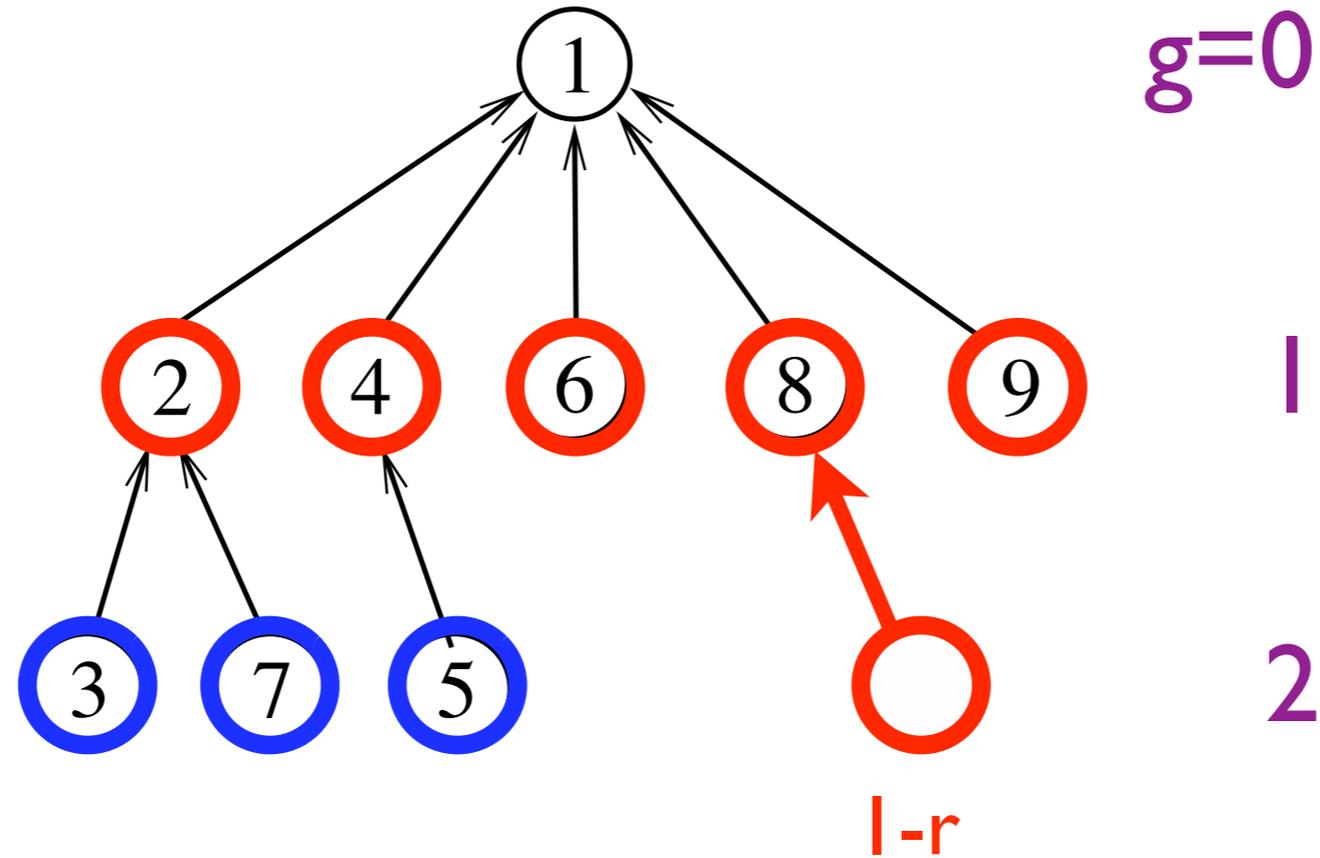
Genealogy (for $A_k=k$) *use redirection!*



Rate equation:

$$\frac{dL_g}{dN} = \frac{(1-r)L_{g-1} + rL_g}{N}$$

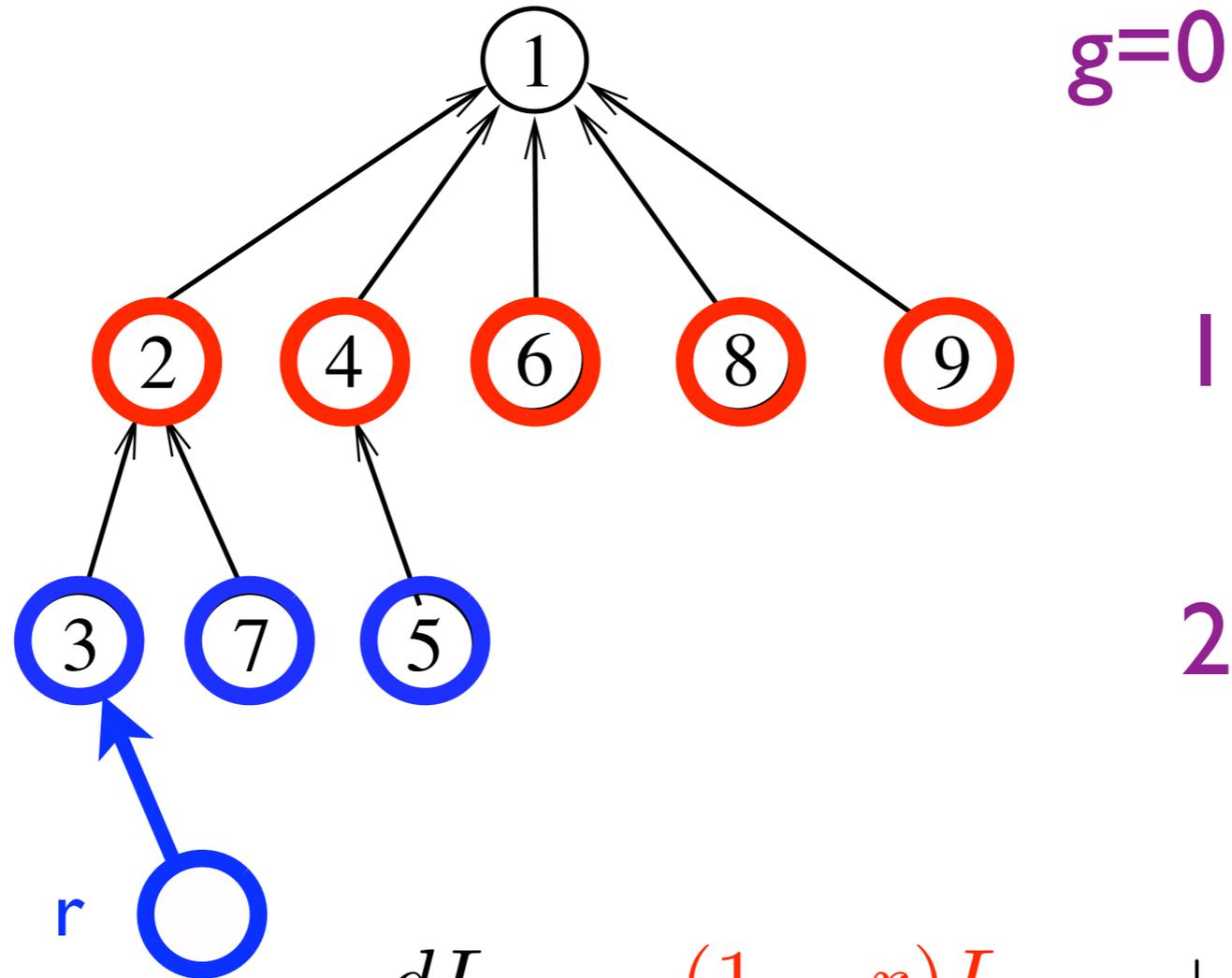
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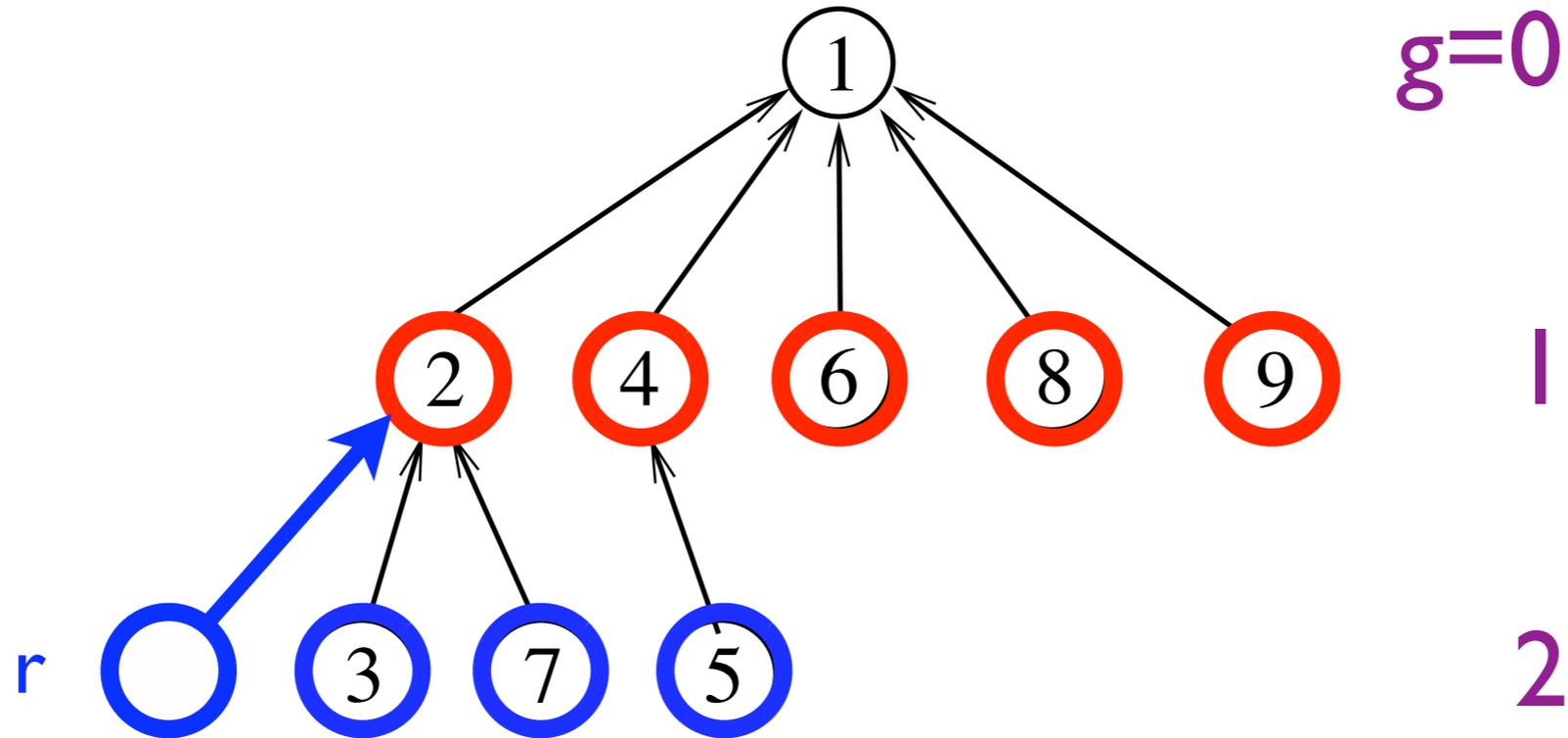
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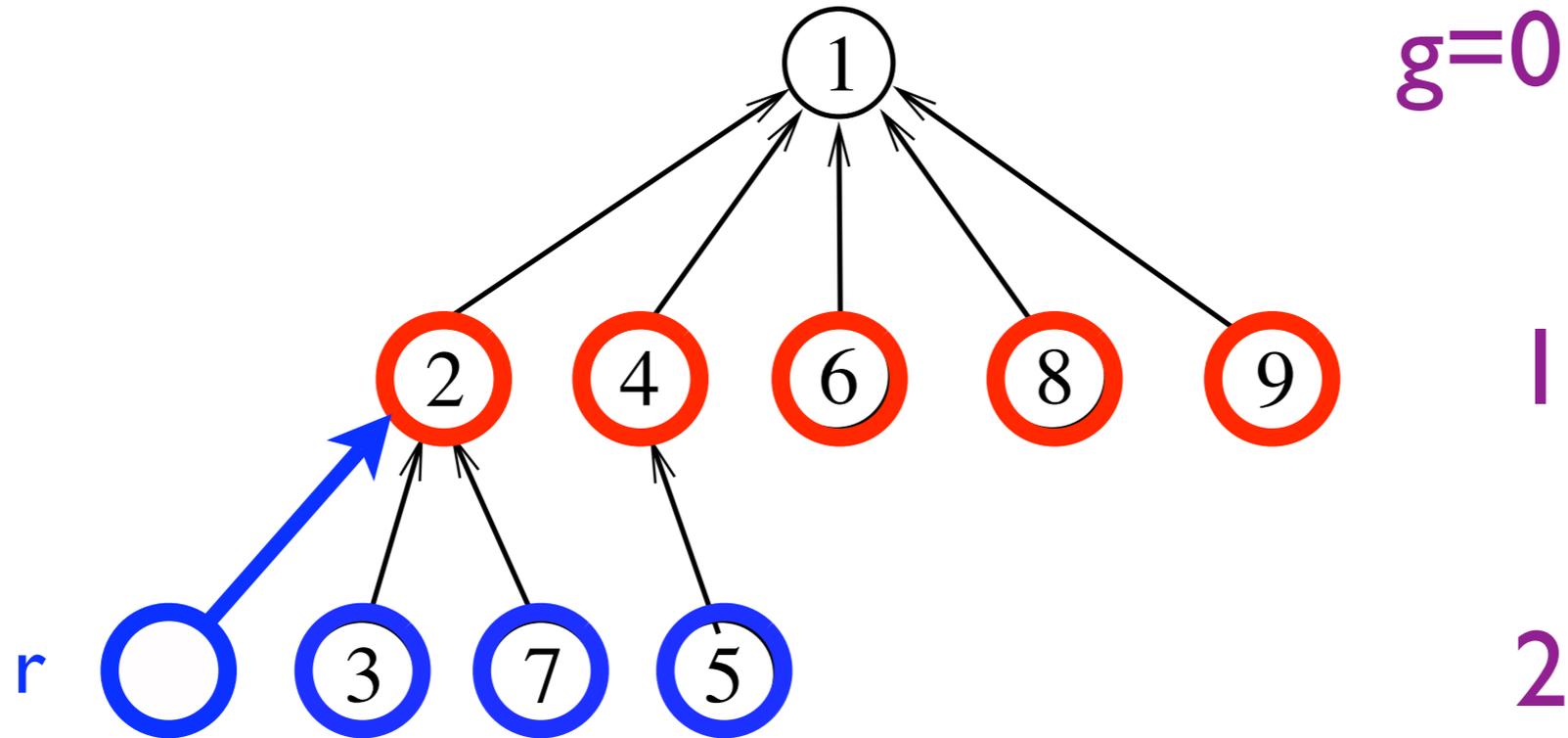


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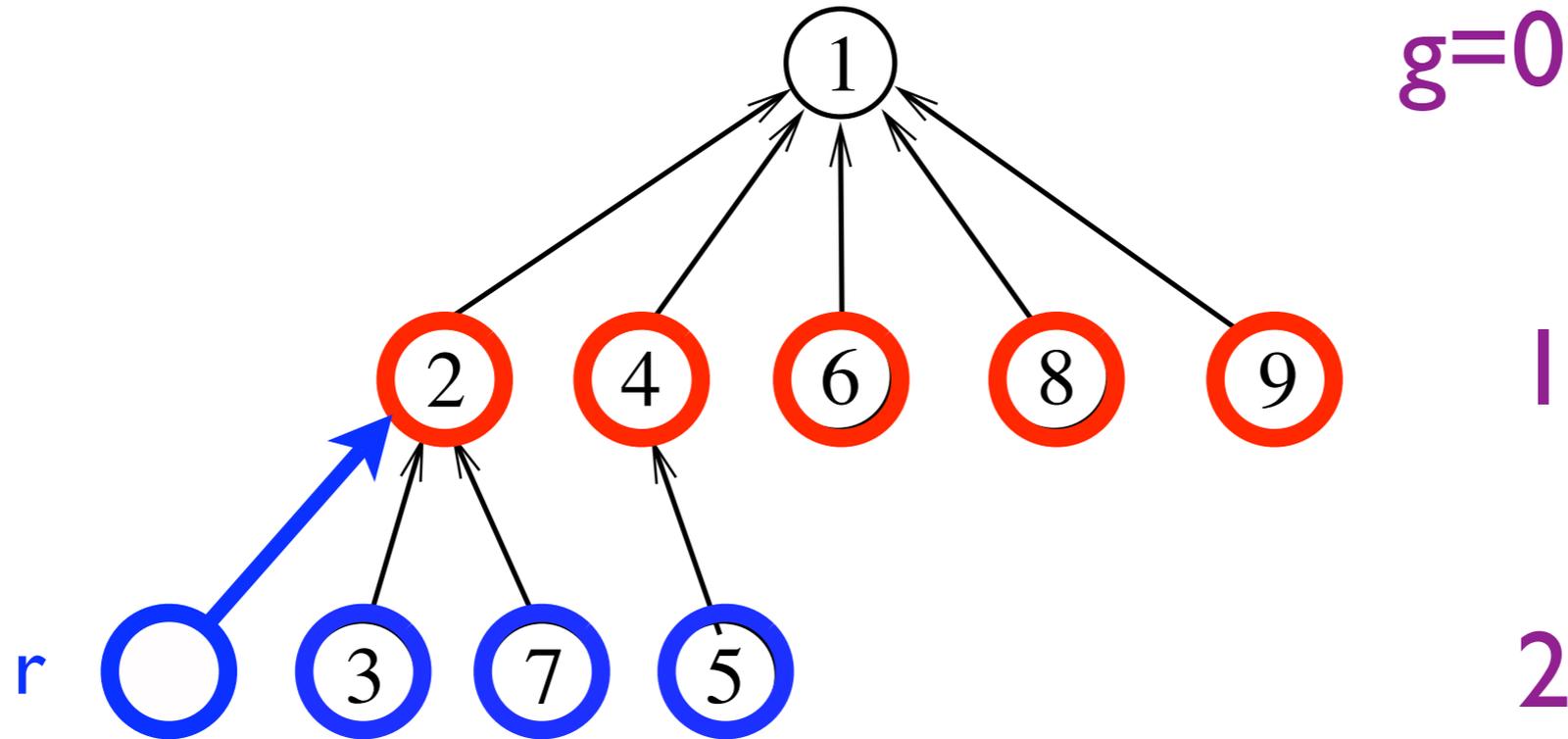


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Outlook

The master equation is a powerful tool to analyze incrementally growing, complex networks.

Wide range of degree distributions arise by preferential attachment; *non-universal power law for linear preferential attachment.*

Topological features and more general models are treatable by the ME approach.

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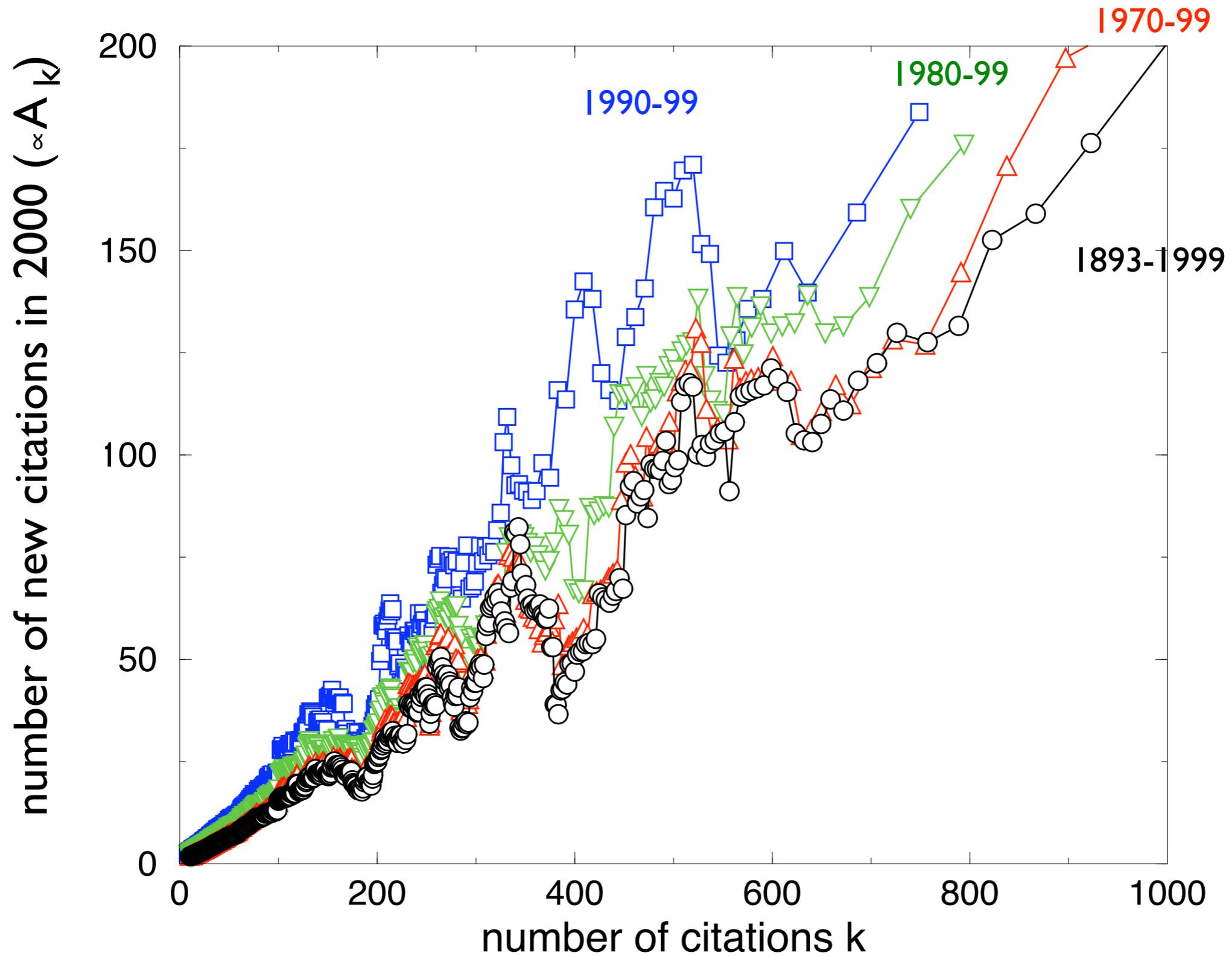
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Distribution of Citations of Physical Review



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